Formal Methods for Security

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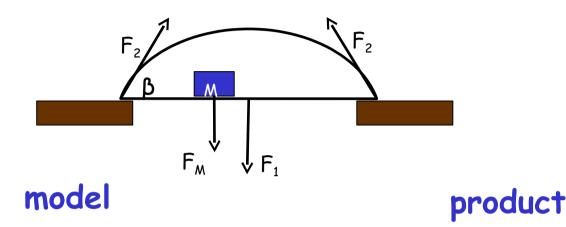
Overview

- Formal methods
 - in general and for security
- Case studies: formal methods for security protocols
 - to illustrate the different ways in formal methods can be used
 - based on our group's formal & informal investigations





Formal Methods for Structural Engineering





Formal methods involve models of which properties (eg bridge won't collapse) can be specified and verified (modulo modelling & abstraction errors) using some methodology/theory

 $F_1 + F_M = 2 * \sin \beta * F_2$ $F_1 = L * H * \rho$

Starting point for all: specification

• which for a bridge is very simple & unchanged for ages

Formal Methods for Software Engineering??



model??

properties??

Untitled import java.util.*;
import java.text.*; /Rod Bernardson //Date: 02/22/2008 //Chapter 18 Programming Challenge 6 //maglerCards Class Demo lic class DealerCardsDem Roaram aros c static void main(String[] args) ne who's turn to play it is new Dealer(): layer = new CardPlayer(deal); er cplayer = new ComputerPlayer(deal); wCard(); .println("Player Points..1: " + makemerision() if (cplayer.getTotalc (cplayer.getTotalCardPoints() <= 21)) pints() > player.getTotalCardPoints() && System.out.println("Computer wins the ("a/a/ lamas System.out.println("\o"): alse if (player.getTotalCardPoints() >
cplayer.getTotalCardPoints() <= 21))</pre> in("Player wins the game! \n\n") else yer.getTotalCardPoints() == .getTotalCardPoints() <= 21)) cplayer.getTotalCardPoints() && System.out.println("Game is a tie! \n\n"); else if (player.getTotalCardPoints() > 21) System.out.println("Game Over - Computer Wins and

specs incl. functional requirements security requirements product, ie code

From specs to code





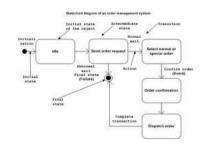


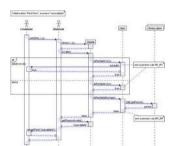




code itself is also possible formal model !

Control Name	190/1EC 27002:2007	Control Function					
		Deter	Avoid	Pro-	De- Sect	React	Re-
Security policy	5	10000					
Information Security Policy	5.1						
Information security policy document	5.1.1	1	1	1		1	1
Review of the information security policy	5.1.2	1	1	1		1	1
Organisation and information security							
Internal organisation	6.1						
Management commitment to information security	6.1.1	1	1	1		1	1
Information security co-ordination	6.1.2		1	1		1	1
Allocation of information security responsibilities	6.1.3		1	~	×	1	×
Authorisation process for information processing facilities	6.1.4		1				
Confidentiality agreements	6.1.5		1	~			
Contact with authorities	6.1.6		1			1	
Contact with special interest groups	6.1.7		1	~	1		
Independent review of information security	6.1.8		1	1	1	1	1
External parties	6.2						
Identification of risks related to external parties	6.2.1	1	1	~			
Addressing security when dealing with customers	6.2.2	1	1	1			
Addressing security in third party agreements	6.2.3	1	1	1	1	1	1
Asset management	7						
Responsibility for assets	7.1						
Inventory of assets	7.1.1		1	~			×
Ovmership of assets	7.1.2		1		× .	1	1
Acceptable use of assets	7.1.3		1	1			





candidate formal models?

Formal methods at different levels

- Formal methods for *programming languages*, eg
 - type system to rule out buffer overflows
 - static analysis to detect XSS vulnerabilities
- Formal methods for abstract algorithms & protocols, eg
 - prove that your shortest path algorithm is functionally correct
 - prove that HTTPS is secure
- Formal methods for programs, eg
 - prove that a program never throws a NullPointerException
 - prove that a program correctly implements HTTPS

security vs correctness

- A program is correct if it does what it should do
 - ie. *presence* of the *right* behaviour, under normal circumstances
- A program is secure if it is does not do what it should not do
 - ie. *absence* of *insecure* behaviour, under *any* circumstances
 - easy to overlook, and hard to check (eg by testing)
- A program also has to be correct for it to be secure?

Good news: some (generic) security requirements are independent of any detailed functional spec (eg absence of integer overflows)

Bad news: security requirements may be hard to pin down (what does it mean for a system to be secure?) Case studies: formal methods for (implementations of) security protocols

Security protocols

- Why security protocols?
 - they are security-critical components in systems
 - eg HTTPS, EMV (Chip & PIN), electronic passports, ...
 - they are small but complex
 - they have clear security objectives

Note:

- forget about crypto, it's the protocols that matter!
- we can study the abstract protocols, or their concrete implementations

Potential problems in security protocols

- 1. using insecure cryptographic primitives (eg. Oyster card)
- 2. using default keys (eg. lots of systems)
- 3. using an buggy protocol. Security protocols are tricky to get right!
- 4. using an buggy implementation. Software bugs can break
 - a) correctness Easy to detect, since the implementation won't work
 - b) security, by erroneously accepting or crashing on
 - incorrect (malformed) message or
 - incorrect order of messages.

This is harder to detect, since the implementation will work

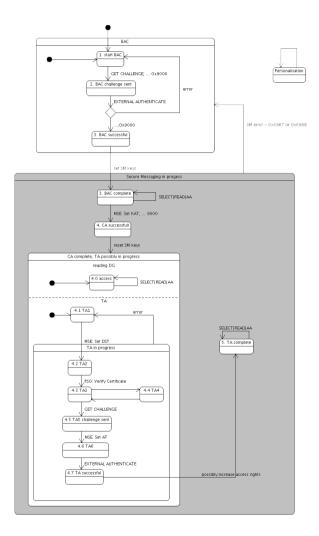
Some example formal models for security protocols

Alice-Bob notation

- 1. A -> B: start session
- 2. B -> A: ok
- 3. A \rightarrow B: Nonce_A
- 4. B -> A: encrypt_{KEY}(Nonce_A)
- 5. A -> B: ...
- 6. B -> A: ...

Such (partial) models capture different aspects and hence can be used for different goals and in different ways (see next slides)

state machines / automata



I. Security Protocol Analysis

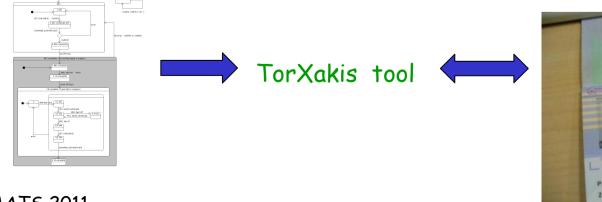
- Given a formal description of the abstract security protocol, eg. in Alice-Bob notation, we can formally analyse some of its properties
 - possible using tool support

Eg next talk by Joeri de Ruiter, and plenty of others.

II. Model based testing

- We automatically test if implementation conforms to the model
 - we feed randomly generated inputs to both model and code, and check if they behave the same
 - the model is used as test oracle
 - possibly also for generating tests & measuring test coverage
- by aggressively testing many (all?) possible sequences we can test for security as well as correctness - "state-based" fuzzing
- Eg we have done this for the electronic passport.

[W.Mostowski et al, FMICS 2009]

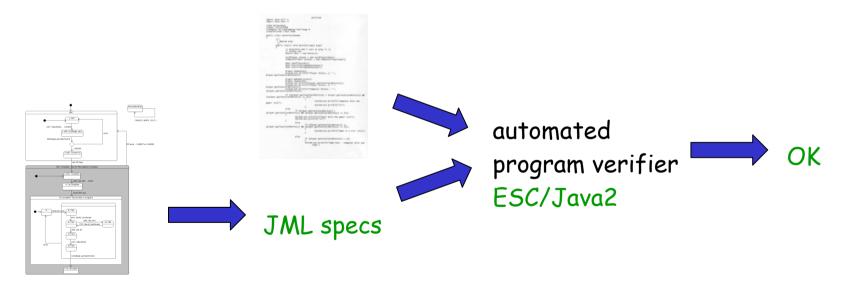




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III. Program verification

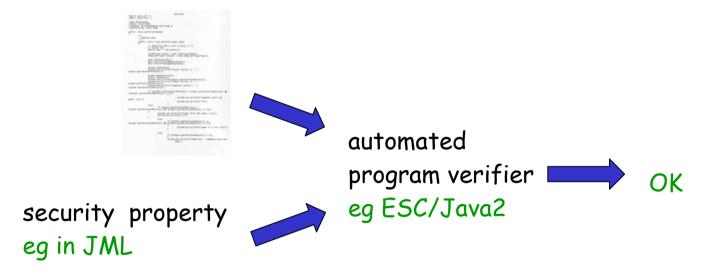
- A more rigorous form of checking compliance of code & model: formal verification (with mathematical proof) that the code conforms to the model
- Eg for a Java implementation of SSH [E.Poll and A.Schubert, WITS 2007]



A formal model can also be used, informally, by a human code reviewer

III. Program verification

• Even without any formal model, we can use formal verification to verify that the code meets some security property



Problem: what do we want to verify anyway?

III. Program specification: what to verify?

Typical easy properties to begin specifying:

(i) important invariants (ii) absence of runtime exceptions plus the additional preconditions and invariants this requires.

public class ElectronicPurse extends javacard.framework.Applet { private int balance; //@ invariant 0 <= balance;</pre>

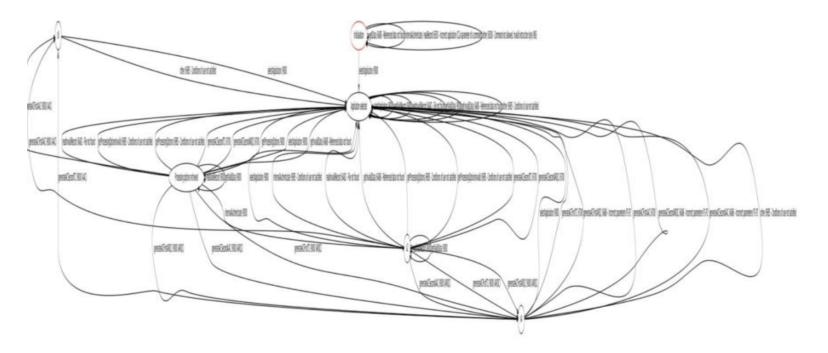
//@ requires buffer != null && 0 <= offset && offset+length <=
buffer.length;</pre>

public static void install (byte[] buffer, short offset, byte length) {

....} Erik Poll, FMATS 2011

IV. Model extraction

• Automated learning techniques can be used (in combination with modelbased testing) to infer an automaton for an implementation's behaviour



Automaton learned from a Dutch EMV bankcard [Fides Aarts et al, ISoLA'10]

Conclusions

- Central challenges
 - does code meet the specs?
 - do specs & code not overlook or introduce security problems?



specs

code

• Formal models & methods can help in different ways

