



Privacy-friendly Electronic Traffic Pricing via Commits



I. ETP Background

ETP, generally

- Replace “flat road tax” by “distance related pricing”
- Pricing may depend on:
 - type of road
 - type of car (esp. emission)
 - time of day
- Aims, apart from fairness,
 - congestion steering/reduction
 - environmental impact reduction



ETP setup via OBU, planned in NL

- Cars get a special box, called **OBU**, for “on-board unit”
- . . . which can at least:
 - determine its own position, via GPS or Galileo
 - communicate with backoffice, via GSM, Wifi, . . .
 - calculate & store data
- Tariff map needed for fee calculation on basis of “trajectory parts”

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Other concerns

- Reliability
- Cost-effectivity (aim in NL: overhead < 10%)
- Ease of use / transparency
- Fraud resistance (e.g. GPS can be manipulated)
- Ease of enforcement
- Ease of dispute resolution
- User acceptance, requiring trust!

Hostile users are to be expected

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Big Question

- Where to store trajectory information:
 - in the back office of the Pricing Authority (PA)
 - in vehicle (i.e. in OBU)
- Architectural decision about information flow
- But also about division of power in society (balance citizen – state)
- **Architecture is Politics!**
(Mitchell Kapor, EFF founder)

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This paper

- Focus on privacy aspects
- Novel solution
(originally due to WdJ; joint elaboration)
- Main ideas will be presented; details may be implemented in several ways
- More general applicability, see later

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II. “Thin” and “Fat”

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ETP via commits

An on-board unit (OBU) is “thin” if:

- it does not itself calculate fees
- Obvious implementation, via centralisation:
 - OBU only computes (somehow) trajectory parts . . .
 - . . . and sends them to the back office of the PA . . .
 - . . . which calculates the fees (and sends bills)
- Easy enforcement via spot checks: take pictures and compare them to PA data

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ETP via commits

Pros and cons of centralisation

- Simple and transparent architecture
- Simple and cheap OBUs
- Extremely privacy-unfriendly
- Central database introduces high risks
 - data compromise may embarrass people (look for politicians who visited prostitute areas)
 - data protection relevant for personal security (e.g. whereabouts of people under threat)
 - single point of failure / bottleneck

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ETP via commits

An OBU is “fat” or “thick” if:

- it calculates fees itself (and passes them on to the PA back office)
- OBU must thus contain tariff map (which must be securely updated, occasionally)
- Spot checks complicated:
 - Two-way communication, while driving by
 - requires integrity & authenticity of OBU
 - requesting most recent trajectory data
 - noticable: generate warning to other drivers

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Pros and cons of “fat”

- Privacy-friendly (via decentralisation)
- Complicated and expensive OBU
- OBU must be trusted
(Successful attack on OBU is catastrophic)
- Complicated spot checks

Essentials

- (I). Fee calculation does not require **identity**
 - Anyone may do it, as long as it can be checked
 - may be a (distributed) service
- (II). Vehicles can **commit** to trajectories, without revealing them: via secure hash!
- (III). Spot checks & fee checks via **revealing** of pre-images

III. Novel solution

Trajectory reporting

- Each OBU calculates, say each minute i , a trajectory part TP_i , from GPS+timing data
- 1440 TPs per day
- At the end it sends the “hash of the day” to the PA back office:

$$h(h^2(TP_1) \parallel \cdots \parallel h^2(TP_{1440}))$$

(h^2 used for fee verification; not essential now)

- SMS size message, that completely fixes trajectories



Trajectory verification

- Assume photo spot check on day d at time t
- PA asks for pre-image of hash of day d :

$$h^2(\text{TP}_1) \parallel \dots \parallel h^2(\text{TP}_{1440})$$

- and asks for pre-images of trajectory parts TP_i around time t — using fixed size of hashes
- Car owner may demand actual photo as evidence — in order to control/limit verification attempts

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Fee calculation

- Assumptions:
 - tariff map is publicly available
 - tariff depends on time of day, type of road, . . .
 - fee for trajectory parts determined by public rules
 - “subfees” should be hidden, as much as possible
- Anyone can calculate fees, eg.
 - one or more (distributed) services
 - own PC, with open source software
 - OBU itself (“fat” version)

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Fee reporting I

- Assume fee reports need to be sent quarterly
- The Pricing Authority (PA) should be able to:
 - check that subfees add up correctly
 - fees of selected trajectory parts are correct
 - . . . in combination with trajectory verification
- In FAST paper realised via **nested hashes**
Here via **homomorphic encryption**

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Fee reporting II

- Assume finite group generator g
Use $f \mapsto g^f$ as hash of fee f
- Multiplication of hashes is sum of fees:

$$g^{f_1} \cdot g^{f_2} = g^{f_1+f_2}$$

- Fees are small numbers that can be tried out: random R needed for blinding in g^{f+R} .

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Fee reporting III; rough version

- A quartely fee report by driver: $\langle \text{Fee}, H, R \rangle$
- where:
 - Fee is total due, over $N = \pm 90$ days
 - H is hash $h(g^{a_1} \parallel \dots \parallel g^{a_N})$ of hashed day fees a_j (including random)
 - R is sum of blindings, with correct sum:
 $g^{\text{Fee}+R} = \prod_i g^{a_i}$
 - $a_j = \sum_{i \leq 1440} \text{Fee}_i + h(\text{TP}_i)$
 - first hash $h(\text{TP}_i)$ of trajectory part is both “blinder” and “binder”

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Fee verification

- Assume PA wants to check time t at day d in fee report $\langle \text{Fee}, H, R \rangle$, after spot check
- PA asks for pre-image day reports g^{a_1}, \dots, g^{a_N} and checks hash H and total sum, via multiplication (and R)
- PA asks for day d all 1440 hashed amounts $g^{\text{Fee}_i + h(\text{TP}_i)}$, which must multiply to g^{a_d}
- PA asks trajectory parts TP_i around t , checks occurrence, and corresponding fees Fee_i .
- This can be fully automated!

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IV. Perspective & conclusions

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Characteristics

- Both privacy-friendly and “thin” possible
- But also “fat”, and more variations (including integration with commercial services)
- Simple spot checking (observation only; should not be noticeable)
- Failure of OBU protection not catastrophic: just increase intensity of spot checks
- Individual remains in control of own data
- Submission of daily hashes compulsory, to enable trajectory (and fee) verification

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Perspective

- General issue “secure & privacy-friendly metering”
- Consumer uses valuables continuously; continuous monitoring by provider
privacy-unfriendly architecture
- Alternative: consumer controls own usage data, but frequently sends hashes; provider monitors occasionally & checks hashes
- Applicable in many more situations
- What do we prefer? Architecture is politics!