

Example of computing a **principal type**:

$$\lambda x^\alpha. \lambda y^\beta. \underbrace{y^\beta (\lambda z^\gamma. \overbrace{y^\beta x^\alpha}^\delta)}_\varepsilon$$

1. Assign type vars to all variables: $x : \alpha, y : \beta, z : \gamma$.
2. Assign type vars to all applicative subterms: $y x : \delta, y(\lambda z. y x) : \varepsilon$.
3. Generate equations between types, necessary for the term to be typable: $\beta = \alpha \rightarrow \delta$ $\beta = (\gamma \rightarrow \delta) \rightarrow \varepsilon$
4. Find a **most general unifier** (a **substitution**) for the type vars that solves the equations: $\alpha := \gamma \rightarrow \delta, \beta := (\gamma \rightarrow \delta) \rightarrow \varepsilon, \delta := \varepsilon$
5. The **principal type** of $\lambda x. \lambda y. y(\lambda z. y x)$ is now

$$(\gamma \rightarrow \varepsilon) \rightarrow ((\gamma \rightarrow \varepsilon) \rightarrow \varepsilon) \rightarrow \varepsilon$$

Exercise: Compute principal types for **S** := $\lambda x. \lambda y. \lambda z. x z (y z)$ and for **M** := $\lambda x. \lambda y. x (y (\lambda z. x z z)) (y (\lambda z. x z z))$.