## Axioms for graph clustering

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Outline

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Conclusion



## Graphs



A symmetric weighted graph (or network) is a pair (V, E) of

- a finite set V of **nodes**, and
- a function  $E: V \times V \rightarrow \mathbb{R}_{\geq 0}$  of edge weights,

such that E(i,j) = E(j,i) for all  $i, j \in V$ .

## Graph clustering



### A clustering C of a graph G = (V, E) is a partition of its nodes.

- Social networks
- Hyperlinks
- Protein interaction networks
- References between mathematical theorems
- Brain parcellation



## Clustering methods

- 1. Clustering function
  - $\hat{C}: \mathsf{Graph} o \mathsf{Clustering}$



- 2. Quality function
  - $Q: \mathsf{Graph} imes \mathsf{Clustering} o \mathbb{R}$
- 3. Quality relation
  - $\cdot \preceq^{G} \cdot \subseteq \mathsf{Clustering} \times \mathsf{Clustering}$

## Clustering methods

- 1. Clustering function  $\hat{\mathcal{C}}: \operatorname{Graph} 
  ightarrow \operatorname{Clustering}$
- 2. Quality function

 $Q:\mathsf{Graph}\times\mathsf{Clustering}\to\mathbb{R}$ 

$$Q\left(\begin{array}{c} b \\ c \\ c \\ \end{array}\right) = 0.1234$$

3. Quality relation

 $\cdot \preceq^{G} \cdot \subseteq \mathsf{Clustering} \times \mathsf{Clustering}$ 

## Clustering methods

- 1. Clustering function  $\hat{\mathcal{C}}: \operatorname{Graph} 
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- 2. Quality function  $Q: \mathsf{Graph} \times \mathsf{Clustering} \to \mathbb{R}$
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- Graph clustering is NP hard.
- Top down:

find best cut and repeat

• Bottom up:

group nodes together

• Simulated annealing



- V.D. Blondel, JL. Guillaume, R. Lambiotte, E. Lefebvre Fast unfolding of communities in large networks J. Stat. Mech. 2008
- Best graph clustering method in surveys.
- Method:
  - Move nodes into neighboring clusters to improve quality.
  - 2 Repeat until local maximum.
  - 8 Now cluster the clusters.





























## Some quality functions

- Connected components
- Total weight of within cluster edges  $Q(G, C) = \sum_{c \in C} w_c$
- Modularity

$$Q(G,C) = \sum_{c \in C} (w_c/v_V - (v_c/v_V)^2)$$

• Many more  $Q(G, C) = \sum_{c \in C} -w_c \log(v_c/v_V)$ 



### Families of quality functions

- Connected components with threshold
- Total weight of within cluster edges with penalty  $Q(G, C) = \sum_{c \in C} w_c - \alpha |C|$
- Modularity

$$Q_{\mathsf{RB}}^{\gamma}(G,C) = \sum_{c \in C} (w_c/v_V - \gamma(v_c/v_V)^2)$$

• Many more  $Q(G, C) = \sum_{c \in C} -w_c \log(v_c/\alpha)$ 



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- Which of these quality functions are good?
- There is no good definition of clustering.
- Can we formalize our intuition?
- Previous work is about distance based clustering (hierarchical clustering, K-means, etc.)
- What about graphs?









A quality function Q is scale invariant if

- for all graphs G = (V, E),
- all constants  $\alpha > 0$ ,

 $Q(G, C_1) \ge Q(G, C_2)$  if and only if  $Q(\alpha G, C_1) \ge Q(\alpha G, C_2)$ .

Intuition: Only the edge weights should matter.



Intuition: Only the edge weights should matter.

A quality function Q is **permutation invariant** if

$$Q(G,C) = Q(f(G),f(C)).$$

for all

- graphs G = (V, E) and
- all isomorphisms  $f: V \to V'$ ,

where f is extended to graphs and clusterings in the obvious way.

### Axiom 3: Richness

Intuition:

• All clusterings must be possible.

So,

- no trivial quality functions.
- no fixed number of clusters.

A quality function Q is **rich** if

- for all sets V and
- all partitions  $C^*$  of V,

there is

- a graph G = (V, E)
- such that  $C^*$  is the optimal clustering of G.

Intuition: Adding edges inside a cluster or removing edges between clusters does not make the clustering worse.



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Let

- G = (V, E) and G' = (V, E') be graphs, and
- C be a clustering of G and G'.

Then G' is a *C*-consistent improvement of G if

- $E'(i,j) \ge E(i,j)$  for all  $i \sim_C j$  and
- $E'(i,j) \leq E(i,j)$  for all  $i \not\sim_C j$ .

Intuition: Adding edges inside a cluster or removing edges between clusters does not make the clustering worse.

# A quality function Q is **monotonic** if $Q(G', C) \ge Q(G, C)$ .

for all

- graphs G,
- all clusterings C of G and
- all C-consistent improvements G' of G.







Two graphs  $G_1$  and  $G_2$  agree on the neighborhood of  $V_a \subseteq V_1 \cap V_2$  if  $E_1(i,j) = E_2(i,j)$  for all  $i \in V_a$ ,  $j \in V_1 \cap V_2$ , and  $E_1(i,j) = 0$  for all  $i \in V_a$ ,  $j \in V_1 \setminus V_2$ , and  $E_2(i,j) = 0$  for all  $i \in V_a$ ,  $j \in V_2 \setminus V_1$ . So, for nodes/clusters in  $V_a$ , all incident edges are the same.

A quality function Q is **local** if

 for all graphs G<sub>1</sub> = (V<sub>1</sub>, E<sub>1</sub>) and G<sub>2</sub> = (V<sub>2</sub>, E<sub>2</sub>) that agree on a set V<sub>a</sub> and its neighborhood,

• for all clusterings 
$$C_1$$
 of  $V_1 \setminus V_a$ ,  
 $C_2$  of  $V_2 \setminus V_a$  and

 $C_a, D_a$  of  $V_a$ .

 $\begin{array}{ll} \text{if} & Q(\mathit{G}_1, \mathit{C}_a \cup \mathit{C}_1) \geq Q(\mathit{G}_1, \mathit{D}_a \cup \mathit{C}_1) \\ \text{then} & Q(\mathit{G}_2, \mathit{C}_a \cup \mathit{C}_2) \geq Q(\mathit{G}_2, \mathit{D}_a \cup \mathit{C}_2). \end{array} \end{array}$ 

### Theorem (Kleinberg 2002)

There is no clustering function that is permutation invariant, scale invariant, monotonic and rich.

### Theorem (Ackerman, Ben-David 2008)

There is a clustering quality function that is permutation invariant, scale invariant, monotonic and rich.

## Discontinuity is magic

### Theorem

There is a graph clustering function that is scale invariant, permutation invariant, monotonic, rich and local.

 $\hat{C}_{coco}(G) =$  the connected components of G

- Doesn't this contradict Kleinberg's theorem?
- No: edge weight  $0 = \text{distance } \infty$ .
- Connected components are unstable.

## Discontinuity is magic

### Theorem

There is a graph clustering function that is scale invariant, permutation invariant, monotonic, rich and local.

 $\hat{C}_{coco}(G)$  = the connected components of G

 $Q_{\text{coco}}(G, C) = \mathbf{1}[C \text{ are the connected components of } G]$ 

- Doesn't this contradict Kleinberg's theorem?
- No: edge weight  $0 = \text{distance } \infty$ .
- Connected components are unstable.

Intuition:

• Don't allow such unstable quality functions.

A quality function Q is **continuous** if for every graph G and every clustering C of G, a sufficiently small change in the edge weights leads to a small change in the objective value.

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### Intuition:

• Balance within cluster edges against cluster volume.

$$Q_{\text{modularity}}(G, C) = \sum_{c \in C} \left( \frac{w_c}{v_V} - \left( \frac{v_c}{v_V} \right)^2 \right).$$

Where  

$$v_c = \sum_{i \in c} \sum_{j \in V} E(i, j)$$
 volume of cluster  
 $w_c = \sum_{i \in c} \sum_{j \in c} E(i, j)$  within cluster weight

The obvious:

- Modularity is permutation invariant.
- Modularity is scale invariant.
- Modularity is continuous.

The less obvious:

• Modularity is rich.

The bad:

- Modularity is *not* local.
- Modularity is *not* monotonic.



### Modularity is not local



### Modularity is not monotonic



### Idea 1: Fix the scale

$$Q_{M-\text{fixed}}(G,C) = \sum_{c \in C} \left(\frac{w_c}{M} - \left(\frac{v_c}{M}\right)^2\right)$$

### Is it monotonic?

Take  $v_c = w_c + b_c$  (within + between)

$$\frac{\partial Q_{M-\text{fixed}}(G,C)}{\partial w_c} = \frac{1}{M} - \frac{2w_c + 2b_c}{M^2}.$$

This is negative when  $2v_c > M$ , so not monotonic

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$$Q_{M-\text{fixed}}(G,C) = \sum_{c \in C} \left(\frac{w_c}{M} - \left(\frac{v_c}{M}\right)^2\right)$$

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This is negative when  $2v_c > M$ , so not monotonic.

$$Q_{M,\gamma}(G,C) = \sum_{c \in C} \left( \frac{w_c}{M + \gamma v_c} - \left( \frac{v_c}{M + \gamma v_c} \right)^2 \right).$$

Adaptive scale modularity is

- permutation invariant, continuous and local.
- monotonic for all  $M \ge 0$  and  $\gamma \ge 2$ .
- rich for all  $M \ge 0$  and  $\gamma \ge 1$ .
- scale invariant for M = 0.

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## Related quality functions

- When  $\gamma = 0$ , we get fixed scale modularity. Equivalent to other modularity variants.
- When  $\gamma = 0$  and  $M = v_V$ , we get modularity.
- When M = 0 we get

$$Q_{0,\gamma}(G,C) \propto \sum_{c\in C} \left(\frac{w_c}{v_c} - \frac{1}{\gamma}\right),$$

i.e. normalized cut.

• When  $M \to \infty$  we get

$$Q_{\infty,\gamma}(G,C)\propto \sum_{c\in C}w_c,$$

i.e. unnormalized cut.



### Adaptive Scale Modularity behavior

## Take a simple graph: $(w) \xrightarrow{b} (w)$

- Two cliques each with *w* within weight
- Connected by edges with total weight *b*.
- Total volume 2w + 2b.
- What is the behavior of adaptive scale modularity?



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- Graph clustering by optimization.
- 6 axioms for graph clustering quality functions.
- Graph setting allows for locality.
- Modularity is not monotonic.
- Non-monotonicity leads to splitting of cliques.
- Adaptive scale modularity satisfies all axioms.
- Generalizes both modularity and normalized cut.
- Two parameters to control size of clusters.

### Thank you for your attention.

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