

## Epigram - View from the left

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- 1 Presentation and History
  - Presentation
  - History
  - Futur
- 2 Epigram group and projects
  - Creators
  - Contributors
  - Projects using Epigram
- 3 How Epigram works
  - Features
  - Grammar
  - Inductives families
- 4 How to use Epigram
  - What is given to the user
  - How to write an Epigram program
- 5 Demonstration

# Presentation

- fonctionnal programming language with dependent types
- strong type system that express program specifications
- smooth transition from programming to integrated programs and proofs
- exploits propostions as type principle and pattern matching
- created by C.McBride and J.McKinna
- still implemented by the Epigram group in UK
- available on Linux,Windows and Mac OS X

# Timeline

- Late 90s' : Epigram sprout up from the creators minds
- Beginning of the 21<sup>st</sup> century : first release of Epigram
- 20 ?? : Epigram 2

## Why creating Epigram ?

DTs for proofs  $\implies$  DTs for programming



## Creators

### Conor McBride



Lecturer at University of Strathclyde & Alta Systems  
Opponent for Ulf Norell's PhD thesis on Agda  
contact : [conor@strictlypositive.org](mailto:conor@strictlypositive.org)

## Creators

### James McKinna



Assistant Professor at Radboud University  
PhD on proof-carrying code/certified program  
contact : [james.mckinna@cs.ru.nl](mailto:james.mckinna@cs.ru.nl)

## Contributors

### Thorsten Altenkirch



- Reader (Associate Professor) at the School of Computer Science of the University of Nottingham
- Construction of “elegant and efficient” implementations of type theoretic languages  $\Rightarrow$  Epigram

## Contributors

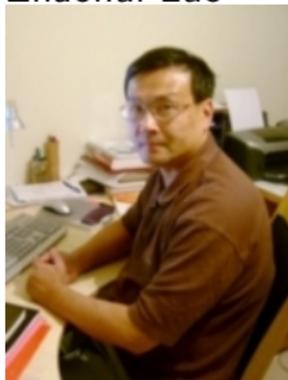
### Healfdene Goguen



- Software engineer at Google, working on large-scale infrastructure
- Co-supervisor of Conor McBride's PhD thesis
- Algorithm for constructor unification for dependent type theories  $\Rightarrow$  Epigram

## Contributors

Zhaohui Luo



- Professor of Computer Science - University of London
- Epigram : Innovative Programming via Inductive Families (J.McKinna, C.McBride, P.Callaghan)

## Contributors

### James Chapman



- Extraordinary Senior Researcher - Tallinn University of Technology.
- Epigram Reloaded : A Standalone Typechecker for ETT (Thorsten Altenkirch, Conor McBride)

## Contributors

Peter Morris



The University of  
**Nottingham**

- Researcher - Department Computer Science University of Nottingham
- Active developer for Epigram

## Contributors

### Edwin Brady



- Researcher - School of Computer Science - University of St Andrews
- Epic : Supercombinator Compiler, intended as a back end for Epigram

## Contributors

Nicolas Oury



- Marie Curie research fellow in the Foundations of Programming Group of the University of Nottingham.
- Involved in the development of Epigram 2

## Projects using Epigram

- Implemented without a module architecture and huge memory footprint  $\Rightarrow$  hard to develop large examples
- Still a “young” programming language
- Active community but relatively small on a very specific topic
- Mainly used for research and publications

# Features

- Pattern matching for types and values based on the data constructors
  - Discriminating tool for DTs
  - Done on the left hand side of intermediate computations, very powerful with DTs
- Propositions  $\iff$  types (Curry-Howard correspondance)
- Programs  $\iff$  proofs
- Deterministic programs
- Types, as well as constraints, induce program structure  $\Rightarrow$  programs as correct by construction

# Grammar

$$\begin{aligned} & lhs \Rightarrow expr \text{ (return instruction)} \\ program & := | lhs \Leftarrow expr \text{ seq } [program] \\ & | lhs | expr [program] \\ decl & := \underline{data} \text{ sig } \underline{where} \text{ sig} \\ & | \underline{let} \text{ sig } program \\ source & := seq [program] \end{aligned}$$

## Inductives families

Class of datatypes aKa Inductives families (Dybjer, 1991)

Inductive family : **So**

$$\frac{\textit{data} \quad b : \mathbf{Bool}}{\mathbf{So} \ b : *}$$

- **So** is a collection of types indexed by a boolean value

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  - and “for free” that  $b$  must be true
- inspecting  $p$  can instantiate  $b \Rightarrow$  any type which depends on either

## What is given to the user

You can found on the Epigram home page :

- All required files to run Epigram (libs + examples)
- IDE (xemacs)
- Literature

# How to write an Epigram program

- Start *xemacs*
- Declare data
- Declare functions and let the editor do (most of) the rest
- Compilation and usage

# Demonstration

- Addition for Nat
- Church-Rosser property for combinatory logic

## Questions

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Sources :

- James McKinna himself ⇒ THANK YOU !
- <http://www.e-pig.org/>
- <http://en.wikipedia.org/wiki/Epigram>
- Conor McBride & James McKinna, The view from the left, Departement of Computer Science - University of Durham