

B Tactics in Fitch style

The green ‘H:’ labels that occur in these descriptions are not part of the way proofs are written in Huth & Ryan, but are necessary for working in ProofWeb. They are the symbolic equivalents (which stay the same throughout the proof process) of the line numbers (which change all the time).

B.1 Structural tactics

exact H						
m	H:	A		m	H:	A
		...				
n		A	→	n		A copy m
insert H B						
	
				n	H:	B
	
n		A	→	$n + 1$		A

B.2 Backward tactics

The tactic names may be prefixed with **b_...** to contrast them to the corresponding forward tactics.

Rules that are not intuitionistically valid are marked with a star. Rules that according to Huth & Ryan are derived rules are marked with a dagger.

<i>conjunction introduction</i>						
con_i						
		
				n		A
	
				$n + 1$		B
n		$A \wedge B$	→	$n + 2$		$A \wedge B$ $\wedge i$ $n, (n + 1)$
<i>conjunction elimination left</i>						
con_e1 B						
		
				n		$A \wedge B$
	
n		A	→	$n + 1$		A $\wedge e_1$ n

conjunction elimination right

con_e2 A

$$\frac{n \quad \dots \quad B \quad \longrightarrow \quad n+1 \quad \dots \quad A \wedge B \quad B \quad \wedge e_2 \quad n}{n+1 \quad B}$$

disjunction introduction left

dis_i1

$$\frac{n \quad \dots \quad A \vee B \quad \longrightarrow \quad n+1 \quad \dots \quad A \quad A \vee B \quad \vee i_1 \quad n}{n+1 \quad A \vee B}$$

disjunction introduction right

dis_i2

$$\frac{n \quad \dots \quad A \vee B \quad \longrightarrow \quad n+1 \quad \dots \quad B \quad A \vee B \quad \vee i_1 \quad n}{n+1 \quad A \vee B}$$

disjunction elimination

dis_e (A ∨ B) H1 H2

$$\frac{\begin{array}{c} \dots \\ n \quad \dots \quad C \quad \longrightarrow \quad n+5 \quad C \quad \vee e \quad n, (n+1) \text{---} (n+2), (n+3) \text{---} (n+4) \end{array}}{\begin{array}{c} \dots \\ n \quad \dots \quad A \vee B \\ n+1 \quad \boxed{\text{H1: } A \quad \text{assumption}} \\ \dots \\ n+2 \quad \boxed{C} \\ n+3 \quad \boxed{\text{H2: } B \quad \text{assumption}} \\ \dots \\ n+4 \quad \boxed{C} \end{array}}$$

implication introduction

imp_i H

$$\frac{\begin{array}{c} \dots \\ n \quad \dots \quad A \rightarrow B \quad \longrightarrow \quad n+2 \quad A \rightarrow B \quad \rightarrow i \quad n \text{---} (n+1) \end{array}}{\begin{array}{c} \dots \\ n \quad \dots \\ n+1 \quad \boxed{B} \end{array}}$$

implication elimination

imp_e A

			n	\dots $A \rightarrow B$	
	\dots		$n+1$	\dots A	
n	B	\rightarrow	$n+2$	B	$\rightarrow_e n, (n+1)$

negation introduction

neg_i H

			n	$H: A$ assumption	
	\dots		$n+1$	\dots \perp	
n	$\neg A$	\rightarrow	$n+2$	$\neg A$	$\neg_i n-(n+1)$

negation elimination

neg_e A

			n	\dots $\neg A$	
	\dots		$n+1$	\dots A	
n	\perp	\rightarrow	$n+2$	\perp	$\neg_e n, (n+1)$

falsum elimination

fls_e

			n	\dots \perp	
n	A	\rightarrow	$n+1$	A	$\perp_e n$

truth introduction

tru_i

n	\dots \top	\rightarrow	n	\top	$\top_i n$
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double negation introduction[†]

negneg_i

			n	\dots A	
n	\dots $\neg\neg A$	\rightarrow	$n+1$	$\neg\neg A$	$\neg\neg_i n$

*double negation elimination**

negneg_e

	...		n	...
n	A	\rightarrow	$n+1$	$\neg\neg A$
				A
				$\neg\neg e\ n$

law of excluded middle†*

LEM

	...		n	
n	$A \vee \neg A$	\rightarrow	n	$A \vee \neg A$
				LEM

proof by contradiction†*

PBC H

			n	H: $\neg A$	assumption
				...	
			$n+1$	\perp	
n	A	\rightarrow	$n+2$	A	PBC $n-(n+1)$

modus tollens†

MT B

			n	...	
				$A \rightarrow B$	
				...	
			$n+1$	$\neg B$	
n	$\neg A$	\rightarrow	$n+2$	$\neg A$	MT $n, (n+1)$

universal introduction

all_i y

			n	y	
				...	
				$A[y/x]$	
n	$\forall x A$	\rightarrow	$n+1$	$\forall x A$	$\forall i\ n-n$

universal elimination

all_e (all x, A)

			n	...	
				$\forall x A$	
n	$A[t/x]$	\rightarrow	$n+1$	$A[t/x]$	$\forall e\ n$

existential introduction

exi_i t

$$\begin{array}{ccc}
 \dots & & \dots \\
 n & \exists x A & \longrightarrow & n & A[t/x] \\
 & & & n+1 & \exists x A & \exists i n
 \end{array}$$

existential elimination

exi_e ($\text{exi } x, A$) y H

$$\begin{array}{ccc}
 & & \dots \\
 & & n & \exists x A \\
 & & n+1 & \boxed{\begin{array}{l} y \\ H: A[y/x] \\ \dots \\ B \end{array}} \\
 \dots & & n+2 & B \\
 n & B & \longrightarrow & n+3 & B & \exists e n, (n+1) \text{---}(n+2)
 \end{array}$$

equality introduction

equ_i

$$\begin{array}{ccc}
 \dots & & \dots \\
 n & t = t & \longrightarrow & n & t = t & =i
 \end{array}$$

equality elimination, simple version

equ_e ($t_1 = t_2$)

$$\begin{array}{ccc}
 & & \dots \\
 & & n & t_1 = t_2 \\
 & & \dots \\
 \dots & & n+1 & A[t_1/x] \\
 n & A[t_2/x] & \longrightarrow & n+2 & A[t_2/x] & =e n, (n+1)
 \end{array}$$

equality elimination, general version (t_2 may occur in A)

equ_e' ($t_1 = t_2$) ($\text{fun } x \Rightarrow A$)

$$\begin{array}{ccc}
 \dots & & \dots \\
 \dots & & n & t_1 = t_2 \\
 & & \dots \\
 \dots & & n+1 & A[t_1/x] \\
 n & A[t_2/x] & \longrightarrow & n+2 & A[t_2/x] & =e n, (n+1)
 \end{array}$$

B.3 Forward versus backward tactics

conjunction introduction
f_con_i H1 H2

m_1 H1 : A	m_1 H1 : A
m_2 H2 : B	m_2 H2 : B
n $A \wedge B$	n $A \wedge B$ $\wedge_i m_1, m_2$

conjunction elimination left
f_con_e1 H

m H : $A \wedge B$	m H : $A \wedge B$
n A	n A $\wedge_{e1} m$

conjunction elimination right
f_con_e2 H

m H : $A \wedge B$	m H : $A \wedge B$
n B	n B $\wedge_{e2} m$

disjunction introduction left
f_dis_i1 H

m H : A	m H : A
n $A \vee B$	n $A \vee B$ $\vee_{i1} m$

disjunction introduction right
f_dis_i2 H

m H : B	m H : B
n $A \vee B$	n $A \vee B$ $\vee_{i2} m$

disjunction elimination

f_dis_e H H1 H2

m **H**: $A \vee B$

m **H**: $A \vee B$

n **H1**: A assumption
 \dots
 $n+1$ C

$n+2$ **H2**: B assumption
 \dots
 $n+3$ C

\dots
 n C

\longrightarrow

$n+4$ C $\vee e$ $m, n-(n+1), (n+2)-(n+3)$

implication elimination

f_imp_e H1 H2

m_1 **H1**: $A \rightarrow B$

m_1 **H1**: $A \rightarrow B$

m_2 **H2**: A

m_2 **H2**: A

\dots
 n B

\longrightarrow

n B $\rightarrow e$ m_1, m_2

negation elimination

f_neg_e H1 H2

m_1 **H1**: $\neg A$

m_1 **H1**: $\neg A$

m_2 **H2**: A

m_2 **H2**: A

\dots
 n \perp

\longrightarrow

n \perp $\neg e$ m_1, m_2

falsum elimination

f_fls_e H

m **H**: \perp

m **H**: \perp

\dots
 n A

\longrightarrow

n A $\perp e$ m

truth introduction
f_tru_i

n	\dots \top	\longrightarrow	n	\top	$\top i n$
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double negation introduction[†]
f_negneg_i H

m	H: A		m	H: A	
n	\dots $\neg\neg A$	\longrightarrow	n	$\neg\neg A$	$\neg\neg i m$

*double negation elimination**
f_negneg_e H

m	H: $\neg\neg A$		m	H: $\neg\neg A$	
n	\dots A	\longrightarrow	n	A	$\neg\neg e m$

law of excluded middle[†]*
f_LEM

n	\dots $A \vee \neg A$	\longrightarrow	n	$A \vee \neg A$	LEM
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modus tollens[†]
f_MT H1 H2

m_1	H1: $A \rightarrow B$		m_1	H1: $A \rightarrow B$	
m_2	H2: $\neg B$		m_2	H2: $\neg B$	
n	\dots $\neg A$	\longrightarrow	n	$\neg A$	MT m_1, m_2

universal elimination
f_all_e H

m	H: $\forall x A$		m	H: $\forall x A$	
n	\dots $A[t/x]$	\longrightarrow	n	$A[t/x]$	$\forall e m$

existential introduction

f_exi_i H

$$\begin{array}{ccc} m & \mathbf{H}: A[t/x] & \\ \dots & & \\ n & \exists x A & \longrightarrow \\ & & m \quad \mathbf{H}: A[t/x] \\ & & n \quad \exists x A \quad \exists i m \end{array}$$

existential elimination

f_exi_e H y H1

$$\begin{array}{ccc} m & \mathbf{H}: \exists x A & \\ & & m \quad \mathbf{H}: \exists x A \\ & & \boxed{\begin{array}{l} y \\ n \quad \mathbf{H1}: A[y/x] \\ \dots \\ n+1 \quad B \end{array}} \\ n & B & \longrightarrow \\ & & n+2 \quad B \quad \exists e m, n-(n+1) \end{array}$$

equality introduction

f_equ_i

$$\begin{array}{ccc} & \dots & \\ n & t = t & \longrightarrow \\ & & n \quad t = t \quad =i \end{array}$$

equality elimination

f_equ_e H1 H2

$$\begin{array}{ccc} m_1 & \mathbf{H1}: t_1 = t_2 & \\ m_2 & \mathbf{H2}: A[t_1/x] & \\ & \dots & \\ n & A[t_2/x] & \longrightarrow \\ & & m_1 \quad \mathbf{H1}: t_1 = t_2 \\ & & m_2 \quad \mathbf{H2}: A[t_1/x] \\ & & n \quad A[t_2/x] \quad =e m_1, m_2 \end{array}$$
