

Proving with Computer Assistance, 2IMF15

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Exercises on First order dependent type theory, λP , some answers

NB \rightarrow binds strongest.

1. Give a precise derivation of the following judgment. (This is example 3 on the slides of the course.)

$$A : *, P : A \rightarrow *, a : A \vdash (Pa) \rightarrow * : \square$$

(Advise: give the derivation in “flag style”, as it was shown in the lecture.)

ANSWER: We give it completely, using the \rightarrow -formation rule as a degenerate case of the Π -formation rule (if $x \notin \text{FV}(B)$):

$$\frac{\Gamma \vdash A : * \quad \Gamma \vdash B : *}{\Gamma \vdash A \rightarrow B : *} \rightarrow\text{-form} \qquad \frac{\Gamma \vdash A : * \quad \Gamma, x:A \vdash B : *}{\Gamma \vdash \Pi x:A. B : *} \Pi\text{-form}$$

$$\begin{array}{l|l|l} 1 & * : \square & \\ 2 & \left| \frac{}{A : *} \right. & \text{var, 1} \\ 3 & \left| \frac{}{A \rightarrow * : \square} \right. & \rightarrow\text{-form, 2, 1} \\ 4 & \left| \left| \frac{}{P : A \rightarrow *} \right. \right. & \text{var, 3} \\ 5 & \left| \left| \left| \frac{}{a : A} \right. \right. \right. & \text{var, 2} \\ 6 & \left| \left| \left| \frac{}{Pa : *} \right. \right. \right. & \text{app, 4, 5} \\ 7 & \left| \left| \frac{}{Pa \rightarrow * : \square} \right. \right. & \rightarrow\text{-form, 6, 1} \end{array}$$

2. Find a term of the following type and write down the context in which this term is typed. (This is example 5 on the slides of the course.)

$$(\Pi x:A. Px \rightarrow Qx) \rightarrow (\Pi x:A. Px) \rightarrow \Pi x:A. Qx$$

Do this by giving a derivation in “flag style”, where you may omit derivations of the well-formedness of types.

ANSWER: Write σ for $(\Pi x:A. Px \rightarrow Qx) \rightarrow (\Pi x:A. Px) \rightarrow \Pi x:A. Qx$.

1	$A : *$	
2	$P : A \rightarrow *$	
3	$Q : A \rightarrow *$	
4	$h : \Pi x:A.P x \rightarrow Q x$	
5	$g : \Pi x:A.P x$	
6	$x : A$	
7	$h x : P x \rightarrow Q x$	app, 4, 6
8	$g x : P x$	app, 5, 6
9	$h x(g x) : Q x$	app, 7, 8
10	$\lambda x:A.h x(g x) : \Pi x:A.Q x$	λ -rule, 6, 9
11	$\lambda g:\Pi x:A.P x.\lambda x:A.h x(g x) : (\Pi x:A.P x) \rightarrow \Pi x:A.Q x$	λ -rule, 5, 10
12	$\lambda h:\Pi x:A.P x \rightarrow Q x.\lambda g:\Pi x:A.P x.\lambda x:A.h x(g x) : \sigma$	λ -rule, 4, 11

So:

$$A : *, P : A \rightarrow *, Q : A \rightarrow * \vdash \lambda h:\Pi x:A.P x \rightarrow Q x.\lambda g:\Pi x:A.P x.\lambda x:A.h x(g x) : \sigma$$

3. Find a term of the following type and write down the context in which this term is typed.

$$(\Pi x:A.P x \rightarrow \Pi z:A.R x z) \rightarrow (\Pi x:A.P x) \rightarrow \Pi z:A.R z z).$$

(NB. Read this type in the proper way: \rightarrow binds stronger than Π !)

ANSWER: write τ for $(\Pi x:A.P x \rightarrow \Pi z:A.R x z) \rightarrow (\Pi x:A.P x) \rightarrow \Pi z:A.R z z)$

$$A : *, P : A \rightarrow *, R : A \rightarrow A \rightarrow * \vdash \\ \lambda h : \Pi x:A.P x \rightarrow \Pi z:A.R x z.\lambda g : \Pi x:A.P x.\lambda y : A.h y(g y) y : \tau$$

4. Give a term M of type $\Pi x:A.P(f(f x))$ in the context

$$\Gamma := A : *, P : A \rightarrow *, f : A \rightarrow A, g : A \rightarrow A, \\ h : \Pi x:A.P(f x) \rightarrow P(g x), k : \Pi x, y:A.(P x \rightarrow P y) \rightarrow P(f x).$$

Also give a derivation of $\Gamma \vdash M : \Pi x:A.P(f(f x))$ in ‘short form’, so you don’t have to show the well-formedness of the types.

ANSWER (only the term):

$$\lambda x : A.k(f x)(g x)(h x)$$

5. Find a term of the following type and write down the context in which this term is typed.

$$(\Pi x:A.P x \rightarrow Q) \rightarrow (\Pi x:A.P x) \rightarrow Q$$

What is special about your context? It should somehow explicitly state that the type A is not empty. How? Why?

ANSWER:

$$A : *, P : A \rightarrow *, Q : *, \mathbf{a} : \mathbf{A} \vdash \\ \lambda h : \Pi x:A.P x \rightarrow Q. \lambda g : \Pi x:A.P x. h a (g a) : (\Pi x:A.P x \rightarrow Q) \rightarrow (\Pi x:A.P x) \rightarrow Q$$

We need a declaration of a variable $a : A$ in the context, stating that A is not empty. If we don't have a term of type A , we can not construct a term of this type, so if A is just a variable in the context, the only thing we can do is to declare $a : A$ as well.

Note that, if A is the “empty type”, the type $(\Pi x:A.P x \rightarrow Q) \rightarrow (\Pi x:A.P x) \rightarrow Q$, interpreted as a formula states something that is just not true: $\forall x:A.P x \rightarrow Q$ and $\forall x:A.P x$ are both vacuously true if A is empty, but Q need not be.

6. Find a term from the given hypotheses of the following type and write down the context in which this term is typed.

$$\begin{aligned} & \forall x. (P(x) \rightarrow R(x, f(x))), \\ & \forall x, y. (R(x, y) \rightarrow R(y, x)), \\ & \forall x, y. (R(x, y) \rightarrow R(f(y), x)) \quad \vdash \quad \forall x. (P(x) \rightarrow R(f(x), f(x))) \end{aligned}$$

ANSWER: We only give the term, type and context. The derivation you have to add yourself!

$$\begin{aligned} & \text{In context } D : *, f : D \rightarrow D, P : D \rightarrow *, R : D \rightarrow D \rightarrow *, \\ & t : \Pi x:D.P x \rightarrow R x (f x), \\ & s : \Pi x, y:D.R x y \rightarrow R y x, \\ & q : \Pi x, y:D.R x y \rightarrow R (f y) x, \text{ we have} \end{aligned}$$

$$\lambda x:D. \lambda h:P x. q (f x) x (s x (f x) (t x h)) : \Pi x:D.P x \rightarrow R (f x) (f x).$$