

## 2IMF15 Proving with Computer Assistance

### Exercises on the Church-Rosser Property

All exercises are about the Church-Rosser proof from the Takahashi paper that we have discussed at the lecture. We recall the definition of  $\Rightarrow$  using derivation rules:

$$\frac{}{x \Rightarrow x} \text{ (var)} \quad \frac{M \Rightarrow M'}{\lambda x.M \Rightarrow \lambda x.M'} \text{ (\lambda)}$$

$$\frac{M \Rightarrow M' \quad N \Rightarrow N'}{MN \Rightarrow M'N'} \text{ (app)} \quad \frac{M \Rightarrow M' \quad N \Rightarrow N'}{(\lambda x.M)N \Rightarrow M'[N'/x]} \text{ (\beta)}$$

1. Consider the term  $M = (\lambda x y.x x(x y))(\mathbf{II})$ 
  - (a) Give the reduction graph of  $M$ . (You may abbreviate  $\mathbf{II}$  to  $J$  and  $\lambda x y.x x(x y)$  to  $P$ .)
  - (b) Compute  $M^*$  and  $(M^*)^*$ .
  - (c) Prove that  $M \Rightarrow M^*$  and  $M^* \Rightarrow (M^*)^*$  by giving a derivation.
2. In the definition of  $\Rightarrow$ , we change clause  $(\beta)$  into

$$\frac{M \Rightarrow \lambda x.P \quad N \Rightarrow N'}{MN \Rightarrow P[N'/x]}$$

- (a) Give the definition of  $(-)^*$  that goes with this adapted definition of  $\Rightarrow$ .
  - (b) Prove again (with these adapted definitions) that  $M \Rightarrow N$  implies  $N \Rightarrow M^*$ , by doing the inductive step for case  $(\beta)$ .
3. The  $\eta$ -reduction rule is:  $\lambda x.M x \rightarrow_\eta M$ , if  $x \notin \text{FV}(M)$ . In order to prove CR for  $\beta\eta$  we add a clause for  $\eta$ -redexes to the definition of  $\Rightarrow$ :

$$\frac{M \Rightarrow M'}{\lambda x.M x \Rightarrow M'} \quad x \notin \text{FV}(M)$$

- (a) Show that now  $(\lambda y x.y x)\mathbf{I} \Rightarrow \mathbf{I}$ , and show that in the original definition, this is not the case.
- (b) Define  $(-)^*$  for this extension to  $\eta$