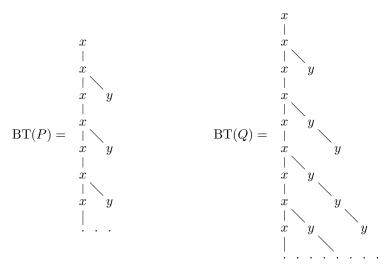
## Semantics and Domain theory

Exercises 13

1. Prove that, for M a closed  $\lambda$ -term, if M has a head-normal-form, then there is a sequence of terms  $P_1, \ldots, P_n$  such that  $M P_1 \ldots P_n =_{\beta} \mathbf{I}$ .

(For closed terms, the reverse implication also holds, so this criterion is equivalent to *having* a *hnf*. This is where the terminology *solvable* comes from.)

- 2. Define  $T := \lambda x.xy(xx)$  and M := TT.
  - (a) Draw the Böhm tree of M.
  - (b) Describe the set of approximations of M,  $\mathcal{A}(M)$ .
- 3. Remember that the **S** combinator is defined as  $\lambda x y z.x z (y z)$ .
  - (a) Draw the Böhm tree of **SSS**.
  - (b) Give the approximations of SSS, that is, describe A(SSS).
- 4. Suppose that the term B satisfies B = x B B. Draw the Böhm tree of B.
- 5. (a) Give a term P that has the Böhm tree given below.
  - (b) (Hard) Give a term Q that has the Böhm tree given below.



6. Let M and N be  $\lambda$ -terms that satisfy the following equations

$$M = \lambda xy.x (M x y) (M x y)$$

$$N = \lambda xy.x (Nxx) (Nxx)$$

Prove that M = N in  $D_A$ .