

RADBOUD UNIVERSITY NIJMEGEN

Science Faculty

Exam **Semantics and Domain Theory** July 2, 2014, 8.30 – 11.30

The maximum number of points per question is given in the margin. (Maximum 100 points in total.) DENS refers to the course notes of Winskell.

NB: you can use all your notes and course material. When dealing with well-known functions and chains that we have seen in the course, you don't have to show (again) that they are monotone/continuous etc.

1. We define the program

$$S := \mathbf{while} \ x > 0 \ \mathbf{do} \ (y := y + x; x := x - 1)$$

Define $w_n : \text{State} \rightarrow \text{State}$ as follows: $w_0(s) = \perp$ and for $n > 0$,

$$w_n(s) := \begin{cases} s & \text{if } s(x) \leq 0 \\ s[x \mapsto 0, y \mapsto s(y) + \frac{s(x)(s(x)+1)}{2}] & \text{if } 0 < s(x) < n \\ \perp & \text{if } n \leq s(x) \end{cases}$$

- (10) (a) Give the functional F that we need in order to compute the denotational semantics of S and show that $F(w_n) = w_{n+1}$

- (10) (b) Give $w_\infty := \sqcup_{n \in \mathbb{N}} w_n$ and show that w_∞ is a fixed-point of F .

- (10) 2. Let f and g be monotone functions from the cpo (D, \sqsubseteq) to (D, \sqsubseteq) . Assume that (i) $f(\perp) = g(\perp)$ and (ii) $f \circ g = g \circ f$.

Prove that $\text{fix } f = \text{fix } g$.

(NB. fix is the well-known Tarski *least-fixed point* of f that we have seen in the course.)

Continue on other side

3. We define the “parallel implication’ function $\text{pimp} : \mathbb{B}_\perp \times \mathbb{B}_\perp \rightarrow \mathbb{B}_\perp$ as follows (the first argument is in horizontal position, the second argument in vertical position):

pimp	\perp	tt	ff
\perp	\perp	\perp	tt
tt	tt	tt	tt
ff	\perp	ff	tt

- (5) (a) Prove that pimp is monotone.
- (5) (b) Prove that pimp is continuous.
- (10) (c) Prove that pimp is not definable in PCF.
- (5) (d) Change one entry in the table of the definition of pimp and show that your adapted function pimp' is definable in PCF.

4. Consider the following PCF terms

$$P_1 := \mathbf{fix} (\mathbf{fn} f : \mathbf{nat} \rightarrow \mathbf{nat} . \mathbf{fn} x : \mathbf{nat} . \mathbf{if} \mathbf{zero} x \mathbf{then} x \mathbf{else} f(\mathbf{pred}(x)))$$

$$P_2 := \mathbf{fn} x : \mathbf{nat} . 0$$

$$P_3 := \mathbf{fn} x : \mathbf{nat} . \mathbf{if} \mathbf{zero} x \mathbf{then} x \mathbf{else} 0$$

- (10) (a) Give a term V such that $P_1(\mathbf{succ}(0)) \Downarrow_{\mathbf{nat}} V$ and give a derivation of $P_1(\mathbf{succ}(0)) \Downarrow_{\mathbf{nat}} V$.
- (10) (b) Give $[P_1]$ and show how you have obtained that answer. (You don’t have to give a full computation.)
- (10) (c) For which i, j do we have $[P_i] \neq [P_j]$?
- (8) 5. Suppose that the λ -terms M_1, M_2 and M_3 satisfy

$$M_1 =_\beta \lambda x . x M_2 (\lambda y . x M_1)$$

$$M_2 =_\beta \lambda x . x M_1 (\lambda y . y M_3)$$

$$M_3 =_\beta \lambda x . x M_2 (\lambda y . x M_3)$$

For which i, j (with $i \neq j$) is it the case that $D_A \models M_i = M_j$? Prove your answer.

- (7) 6. Let M and N be λ -terms and assume that M does not have a head-normal form, while N does have a head-normal form.
Show that in the model D_A , $[\lambda x . x M] \subseteq [\lambda x . x N]$.

END