

# Semantics and Domain theory

## Exercises 11

Recall the following terms in untyped  $\lambda$ -calculus

- $c_n$  denotes the  $n$ -th Church numeral, so in particular  $c_0 = \lambda f x.x$ ,  $c_1 = \lambda f x.f x$  and in general  $c_n = \lambda f x.f^n(x)$ .
  - $\mathbf{K} = \lambda x y.x$ ,  $\mathbf{I} = \lambda x.x$ ,  $\Omega = (\lambda x.x x)(\lambda x.x x)$ .
1. (a) Prove that the equation  $c_0 = c_1$  is inconsistent in untyped  $\lambda$  calculus. (That is: show that, if you assume  $c_0 = c_1$ , then you can prove  $M = N$  for all terms  $M, N$ .)  
(b) Prove that  $c_0 = c_{n+1}$  is inconsistent for any  $n \in \mathbb{N}$ .  
(c) Prove that  $c_n = c_m$  is inconsistent for  $n, m \in \mathbb{N}$  with  $n \neq m$ .

2. The applicative structure  $(M, \cdot)$ , with  $M = \mathbb{N}$  (the natural numbers) and  $\cdot = *$  (multiplication) cannot be made into a (consistent) model of the untyped  $\lambda$ -calculus. We prove this in the following steps:

Consider  $(\mathbb{N}, *)$  and assume that there is an interpretation  $\llbracket - \rrbracket$  satisfying the definition that we have given in the lecture (see Definition 59 of Berline). Now:

- (a) Show that  $\llbracket \mathbf{K I} \rrbracket = \llbracket \mathbf{I K} \rrbracket$ .
- (b) Conclude that  $d = e$  for all  $d, e \in \mathbb{N}$ .

So: all elements are equal in the model.

3. Prove that the theory that equates all  $\lambda$ -terms that don't have a normal form is inconsistent by showing that the following equation is inconsistent in untyped  $\lambda$  calculus:

$$\lambda x y.x y \Omega = \lambda x y.y x \Omega.$$