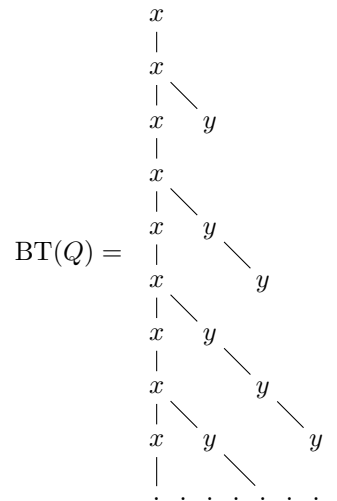
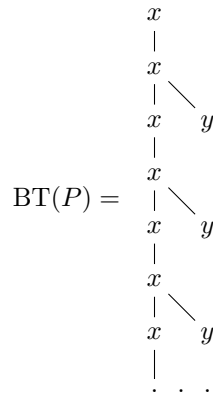


Semantics and Domain theory

Exercises 13

1. Prove that, for M a closed λ -term, if M has a head-normal-form, then there is a sequence of terms P_1, \dots, P_n such that $M P_1 \dots P_n =_{\beta} \mathbf{I}$.
(For closed terms, the reverse implication also holds, so this criterion is equivalent to *having a hnf*. This is where the terminology *solvable* comes from.)
2. Define $T := \lambda x.x y (x x)$ and $M := T T$.
 - (a) Draw the Böhm tree of M .
 - (b) Describe the set of approximations of M , $\mathcal{A}(M)$.
3. Remember that the **S** combinator is defined as $\lambda x y z.x z (y z)$.
 - (a) Draw the Böhm tree of **SSS**.
 - (b) Give the approximations of **SSS**, that is, describe $\mathcal{A}(\mathbf{SSS})$.
4. Suppose that the term B satisfies $B = x B B$. Draw the Böhm tree of B .
5. (a) Give a term P that has the Böhm tree given below.
(b) (Hard) Give a term Q that has the Böhm tree given below.



6. Let M and N be λ -terms that satisfy the following equations

$$M = \lambda x y.x (M x y) (M x y)$$

$$N = \lambda x y.x (N x x) (N x x)$$

Prove that $M = N$ in D_A .