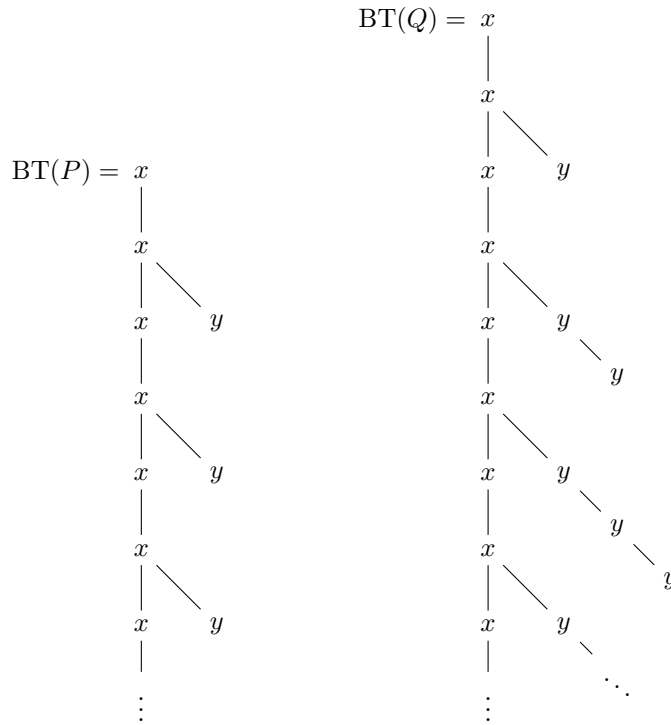


# Semantics and Domain theory

## Exercises 13

1. Prove that, for  $M$  a closed  $\lambda$ -term, if  $M$  has a head-normal-form, then there is a sequence of terms  $P_1, \dots, P_n$  such that  $M P_1 \dots P_n =_{\beta} \mathbf{I}$ .  
(For closed terms, the reverse implication also holds, so this criterion is equivalent to *having a hnf*. This is where the terminology *solvable* comes from.)
2. Define  $T := \lambda x.x y (x x)$  and  $M := T T$ .
  - (a) Draw the Böhm tree of  $M$ .
  - (b) Describe the set of approximations of  $M$ ,  $\mathcal{A}(M)$ .
3. Remember that the **S** combinator is defined as  $\lambda x y z.x z (y z)$ .
  - (a) Draw the Böhm tree of **SSS**.
  - (b) Give the approximations of **SSS**, that is, describe  $\mathcal{A}(\mathbf{SSS})$ .
4. Suppose that the term  $B$  satisfies  $B = x B B$ . Draw the Böhm tree of  $B$ .
5. (a) Give a term  $P$  that has the Böhm tree given below.  
(b) (Hard) Give a term  $Q$  that has the Böhm tree given below.



6. Let  $M$  and  $N$  be  $\lambda$ -terms that satisfy the following equations

$$M = \lambda x y.x (M x y) (M x y)$$

$$N = \lambda x y.x (N x x) (N x x)$$

Prove that  $M = N$  in  $D_A$ .