

$[\text{ass}_{\text{ns}}]$	$\langle x := a, s \rangle \rightarrow s[x \mapsto \mathcal{A}[a]]s$
$[\text{skip}_{\text{ns}}]$	$\langle \text{skip}, s \rangle \rightarrow s$
$[\text{comp}_{\text{ns}}]$	$\frac{\langle S_1, s \rangle \rightarrow s', \langle S_2, s' \rangle \rightarrow s''}{\langle S_1; S_2, s \rangle \rightarrow s''}$
$[\text{if}_{\text{ns}}^{\text{tt}}]$	$\frac{\langle S_1, s \rangle \rightarrow s'}{\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \rightarrow s'} \quad \text{if } \mathcal{B}[b]s = \text{tt}$
$[\text{if}_{\text{ns}}^{\text{ff}}]$	$\frac{\langle S_2, s \rangle \rightarrow s'}{\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \rightarrow s'} \quad \text{if } \mathcal{B}[b]s = \text{ff}$
$[\text{while}_{\text{ns}}^{\text{tt}}]$	$\frac{\langle S, s \rangle \rightarrow s', \langle \text{while } b \text{ do } S, s' \rangle \rightarrow s''}{\langle \text{while } b \text{ do } S, s \rangle \rightarrow s''} \quad \text{if } \mathcal{B}[b]s = \text{tt}$
$[\text{while}_{\text{ns}}^{\text{ff}}]$	$\langle \text{while } b \text{ do } S, s \rangle \rightarrow s \quad \text{if } \mathcal{B}[b]s = \text{ff}$

 Table 2.1: Natural semantics for **While**

$[\text{ass}_{\text{sos}}]$	$\langle x := a, s \rangle \Rightarrow s[x \mapsto \mathcal{A}[a]]s$
$[\text{skip}_{\text{sos}}]$	$\langle \text{skip}, s \rangle \Rightarrow s$
$[\text{comp}_{\text{sos}}^1]$	$\frac{\langle S_1, s \rangle \Rightarrow \langle S'_1, s' \rangle}{\langle S_1; S_2, s \rangle \Rightarrow \langle S'_1; S_2, s' \rangle}$
$[\text{comp}_{\text{sos}}^2]$	$\frac{\langle S_1, s \rangle \Rightarrow s'}{\langle S_1; S_2, s \rangle \Rightarrow \langle S_2, s' \rangle}$
$[\text{if}_{\text{sos}}^{\text{tt}}]$	$\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \Rightarrow \langle S_1, s \rangle \quad \text{if } \mathcal{B}[b]s = \text{tt}$
$[\text{if}_{\text{sos}}^{\text{ff}}]$	$\langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \Rightarrow \langle S_2, s \rangle \quad \text{if } \mathcal{B}[b]s = \text{ff}$
$[\text{while}_{\text{sos}}]$	$\langle \text{while } b \text{ do } S, s \rangle \Rightarrow$ $\langle \text{if } b \text{ then } (S; \text{while } b \text{ do } S) \text{ else skip}, s \rangle$

 Table 2.2: Structural operational semantics for **While**

$[\text{ass}_{\text{p}}]$	$\{ P[x \mapsto \mathcal{A}[a]] \} x := a \{ P \}$
$[\text{skip}_{\text{p}}]$	$\{ P \} \text{skip} \{ P \}$
$[\text{comp}_{\text{p}}]$	$\frac{\{ P \} S_1 \{ Q \}, \{ Q \} S_2 \{ R \}}{\{ P \} S_1; S_2 \{ R \}}$
$[\text{if}_{\text{p}}]$	$\frac{\{ \mathcal{B}[b] \wedge P \} S_1 \{ Q \}, \{ \neg \mathcal{B}[b] \wedge P \} S_2 \{ Q \}}{\{ P \} \text{if } b \text{ then } S_1 \text{ else } S_2 \{ Q \}}$
$[\text{while}_{\text{p}}]$	$\frac{\{ \mathcal{B}[b] \wedge P \} S \{ P \}}{\{ P \} \text{while } b \text{ do } S \{ \neg \mathcal{B}[b] \wedge P \}}$
$[\text{cons}_{\text{p}}]$	$\frac{\{ P' \} S \{ Q' \}}{\{ P \} S \{ Q \}} \quad \text{if } P \Rightarrow P' \text{ and } Q' \Rightarrow Q$

Table 9.1: Axiomatic system for partial correctness

[block _{ns}]	$\frac{\langle D_V, s \rangle \rightarrow_D s', \langle S, s' \rangle \rightarrow s''}{\langle \text{begin } D_V \ S \ \text{end}, s \rangle \rightarrow s''[\text{DV}(D_V) \mapsto s]}$
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Table 3.1: Natural semantics for statements of **Block**

[var _{ns}]	$\frac{\langle D_V, s[x \mapsto \mathcal{A}[a]]s \rangle \rightarrow_D s'}{\langle \text{var } x := a; D_V, s \rangle \rightarrow_D s'}$
[none _{ns}]	$\langle \varepsilon, s \rangle \rightarrow_D s$

Table 3.2: Natural semantics for variable declarations

[ass _{ns}]	$\text{env}_P \vdash \langle x := a, s \rangle \rightarrow s[x \mapsto \mathcal{A}[a]]s$
[skip _{ns}]	$\text{env}_P \vdash \langle \text{skip}, s \rangle \rightarrow s$
[comp _{ns}]	$\frac{\text{env}_P \vdash \langle S_1, s \rangle \rightarrow s', \text{env}_P \vdash \langle S_2, s' \rangle \rightarrow s''}{\text{env}_P \vdash \langle S_1; S_2, s \rangle \rightarrow s''}$
[iftt _{ns}]	$\frac{\text{env}_P \vdash \langle S_1, s \rangle \rightarrow s'}{\text{env}_P \vdash \langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \rightarrow s'}$ <p style="text-align: center;">if $\mathcal{B}[b]s = \text{tt}$</p>
[ifff _{ns}]	$\frac{\text{env}_P \vdash \langle S_2, s \rangle \rightarrow s'}{\text{env}_P \vdash \langle \text{if } b \text{ then } S_1 \text{ else } S_2, s \rangle \rightarrow s'}$ <p style="text-align: center;">if $\mathcal{B}[b]s = \text{ff}$</p>
[while _{ns} ^{tt}]	$\frac{\text{env}_P \vdash \langle S, s \rangle \rightarrow s', \text{env}_P \vdash \langle \text{while } b \text{ do } S, s' \rangle \rightarrow s''}{\text{env}_P \vdash \langle \text{while } b \text{ do } S, s \rangle \rightarrow s''}$ <p style="text-align: center;">if $\mathcal{B}[b]s = \text{tt}$</p>
[while _{ns} ^{ff}]	$\text{env}_P \vdash \langle \text{while } b \text{ do } S, s \rangle \rightarrow s$ <p style="text-align: center;">if $\mathcal{B}[b]s = \text{ff}$</p>
[block _{ns}]	$\frac{\langle D_V, s \rangle \rightarrow_D s', \text{upd}_P(D_P, \text{env}_P) \vdash \langle S, s' \rangle \rightarrow s''}{\text{env}_P \vdash \langle \text{begin } D_V \ D_P \ S \ \text{end}, s \rangle \rightarrow s''[\text{DV}(D_V) \mapsto s]}$
[call _{ns} ^{rec}]	$\frac{\text{env}_P \vdash \langle S, s \rangle \rightarrow s'}{\text{env}_P \vdash \langle \text{call } p, s \rangle \rightarrow s'} \quad \text{where } \text{env}_P \ p = S$

Table 3.3: Natural semantics for **Proc** with dynamic scope rules

[call _{ns}]	$\frac{\text{env}'_P \vdash \langle S, s \rangle \rightarrow s'}{\text{env}_P \vdash \langle \text{call } p, s \rangle \rightarrow s'}$ <p style="text-align: center;">where $\text{env}_P \ p = (S, \text{env}'_P)$</p>
[call _{ns} ^{rec}]	$\frac{\text{env}'_P[p \mapsto (S, \text{env}'_P)] \vdash \langle S, s \rangle \rightarrow s'}{\text{env}_P \vdash \langle \text{call } p, s \rangle \rightarrow s'}$ <p style="text-align: center;">where $\text{env}_P \ p = (S, \text{env}'_P)$</p>

Table 3.4: Procedure calls in case of mixed scope rules (choose one)