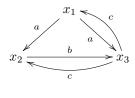
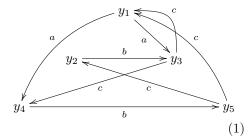
## Exercises Coalgebra for Lecture 7

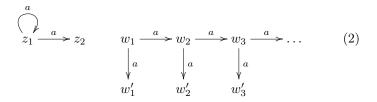
The exercises labeled with (\*) are optional and more advanced.

- 1. For each of the following transition systems and given pairs of states, decide whether they are bisimilar or not (and justify your answer). Are they trace equivalent?
  - (a)  $x_1$  and  $y_1$  in:<sup>1</sup>





(b)  $z_1$  and  $w_1$  in:



(c)  $u_1$  and  $v_1$  in:

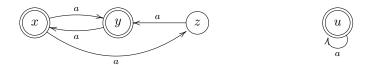
$$u_{1} \xrightarrow{a} u_{2} \xrightarrow{c} u_{3} \qquad v_{2} \qquad v_{4} \qquad (3)$$

$$a \left( \bigvee_{b} b \qquad b \left( \bigvee_{a} a \right) \right) a \qquad v_{1} \longrightarrow v_{3} \longrightarrow v_{5}$$

- 2. Let  $\delta \colon X \to (\mathcal{P}(X))^A$  be a labelled transition system. The *trace semantics* tr:  $X \to 2^{A^*}$  was defined in the lecture.
  - (a) Define a deterministic automaton  $\langle o^{\sharp}, \delta^{\sharp} \rangle \colon \mathcal{P}(X) \to 2 \times (\mathcal{P}(X))^A$  such that the language accepted by a state  $\{x\}$  is given by l(x).
  - (b) (\*) Show that the above construction defines a functor from the category of transition systems (over A) to the category of deterministic automata (over A). What does this say about the relationship between bisimilarity and trace semantics?

 $<sup>^1\</sup>mathrm{Exercise}$  taken from: D. Sangiorgi, Introduction to Bisimulation and Coinduction, Cambridge Universty Press.

3. Prove that x and u in the following non-deterministic automaton accept the same language, using determinisation and bisimulation.



- 4. Show that if two states of an LTS are bisimilar, then they are also trace equivalent.
- 5. In the lecture, we defined the determinisation of a non-deterministic automaton  $\langle o, \delta \rangle \colon X \to 2 \times (\mathcal{P}_f(X))^A$  as the automaton  $\langle o^{\sharp}, \delta^{\sharp} \rangle \colon \mathcal{P}_f(X) \to \mathcal{P}_f(X)$  $2 \times (\mathcal{P}_f(X))^A$  given by

$$o^{\sharp}(S) = \bigvee_{x \in S} o(x) \tag{4}$$

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$$\delta^{\sharp}(S)(a) = \bigcup_{x \in S} \delta(x)(a).$$
(4)

Let beh:  $\mathcal{P}_f(X) \to 2^{A^*}$  be the language semantics of determinisation. Show that beh is a semilattice homomorphism, i.e., that  $beh(\emptyset)(w) = 0$ and  $beh(S \cup T)(w) = beh(S)(w) \vee beh(T)(w)$  for all  $S, T \in \mathcal{P}_f(X)$  and  $w \in A^*$ .

- (a) Show that semilattices and semilattice homomorphisms form a category SL.
  - (b) Does SL have an initial object? And a final object? And products, coproducts?
  - (c) Show that there is a 'forgetful' functor  $U: SL \to Set$ , mapping a semilattice to the underlying set.
  - (d) (\*) Let  $F: \mathsf{Set} \to \mathsf{Set}$  be given by  $F(X) = 2 \times X^A$ . Show that there is a functor  $\overline{F} \colon \mathsf{SL} \to \mathsf{SL}$  such that  $U \circ \overline{F} = F \circ U$  (what does this mean, concretely)?
  - (e) (\*) Let  $G: \mathsf{Set} \to \mathsf{Set}$  be given by  $G(X) = 2 \times (\mathcal{P}_f(X))^A$ . Show that determinisation gives rise to a functor  $\mathsf{CoAlg}(G) \to \mathsf{CoAlg}(\overline{F})$ .