

Reasoning about Quantum Protocols in QPEL

Robin Adams robin@cs.ru.nl

Radboud University

A type theory for denoting quantum programs, and effects on quantum systems. First presented in [1]

Syntax

Types

A type represents a state space for a quantum computer.

I	Unit type
$A \otimes B$	Pair of A and B (possibly entangled)
$A + B$	Either A or B (classical 'or')
qbit	Qubits

Terms

A term represents a quantum computation.

x	Variable
$M \otimes N$	Separated state
let $x \otimes y = M$ in N	Eliminator for $A \otimes B$
inl (M)	Injection $A \rightarrow A + B$
inr (M)	Injection $B \rightarrow A + B$
case M of inl (x) $\rightarrow N$ inr (y) $\rightarrow P$	Eliminator for $A + B$
measure $\phi_1 \rightarrow M_1 \mid \dots \mid \phi_n \rightarrow M_n$	Make measurement then branch
inlr (M, N)	Returns the output of M or N (only if exactly one of M, N terminates)
$ 0\rangle$	New qubit
UM	Apply a quantum gate (U a $2^n \times 2^n$ complex unitary matrix)

Effects

An effect represents a measurement that can be performed.

0	Always false effect
$\phi \oplus \psi$	Sum of ϕ and ψ (only if $\phi \perp \psi$)
ϕ^\perp	Orthocomplement of ϕ
$\phi \otimes \psi$	Tensor product
case M of inl (x) $\rightarrow \phi$ inr (y) $\rightarrow \psi$	Perform one of two measurements depending on M
$ 0\rangle M$	Measure qubit in standard basis

We write 1 for 0^\perp .

Judgements

QPEL has four judgement forms:

$\Gamma \vdash M : A$	M is an algorithm taking inputs Γ , giving output A
$\Gamma \vdash M = N : A$	M and N give the same output to the same input
$\Gamma \vdash \phi$ pred	ϕ is an effect that may be measured on Γ
$\Gamma \vdash \phi \leq \psi$	Probability of ϕ is \leq probability of ψ on all states of Γ

We write $\Gamma \vdash \phi \Leftrightarrow \psi$ for the two judgements $\Gamma \vdash \phi \leq \psi$ and $\Gamma \vdash \psi \leq \phi$.

The effects (measurements) are the *predicates* of our logic.

$\Gamma \vdash \phi \leq \psi$ is the entailment relation.

Quantum Teleportation

The **inlr** construction allows us to reason case-by-case:

$$\begin{aligned}
 F_1(x) &\stackrel{\text{def}}{=} \text{let } x \otimes y \otimes z = \text{prepare } (x) \text{ in } = \text{measure} \begin{cases} 1/4 \rightarrow \text{inl } (x) \\ 3/4 \rightarrow \text{inr } (\diamond) \end{cases} \\
 &\quad \text{measure} \begin{cases} \langle 00 | xy \rightarrow \text{inl } (z) \\ \langle 00 | xy \rangle^\perp \rightarrow \text{inr } (\diamond) \end{cases} \\
 F_2(x) &\stackrel{\text{def}}{=} \text{let } x \otimes y \otimes z = \text{prepare } (x) \text{ in } = \text{measure} \begin{cases} 1/4 \rightarrow \text{inl } (\diamond) \\ 1/4 \rightarrow \text{inr } (\text{inl } (x)) \\ 1/2 \rightarrow \text{inr } (\text{inr } (\diamond)) \end{cases} \\
 &\quad \text{measure} \begin{cases} \langle 00 | xy \rightarrow \text{inl } (\diamond) \\ \langle 01 | xy \rightarrow \text{inr } (\text{inl } (Xz)) \\ \langle 1 | x \rightarrow \text{inr } (\text{inr } (\diamond)) \end{cases} \\
 &\quad \text{THEREFORE} \\
 &\quad (\nabla + \text{id})(\text{inlr } (F_1(x), F_2(x))) \\
 &= \text{let } x \otimes y \otimes z = \text{prepare } (x) \text{ in } = \text{measure} \begin{cases} 1/2 \rightarrow \text{inl } (x) \\ 1/2 \rightarrow \text{inr } (\diamond) \end{cases} \\
 &\quad \text{measure} \begin{cases} \langle 00 | xy \rightarrow \text{inl } (z) \\ \langle 01 | xy \rightarrow \text{inl } (Xz) \\ \langle 1 | x \rightarrow \text{inr } (\diamond) \end{cases}
 \end{aligned}$$

Handling the other two cases similarly, we build up to

$$x : \text{qbit} \vdash \text{teleport } (x) = x : \text{qbit}$$

Semantics

If we omit the type **qbit**, the models of QPEL are *state-and-effect triangles*

$$\begin{array}{ccc}
 \text{EMod}_E^{\text{op}} & \xrightarrow{\quad} & \text{Conv}_E \\
 & \searrow \mathcal{V} & \swarrow S \\
 & &
 \end{array}$$

With **qbit**, the rules are sound and complete for $\mathcal{V} = \text{CStar}_{\text{CPU}}^{\text{op}}$.

Structural Rules

QPEL is an *affine* type system. It satisfies Weakening and Exchange but *not* Contraction.

$$\begin{array}{l}
 \text{(Weakening)} \frac{\Gamma \vdash \mathcal{J} \quad (\Gamma \subseteq \Delta)}{\Delta \vdash \mathcal{J}} \quad \text{(exchange)} \frac{\Gamma, x : A, y : B, \Delta \vdash \mathcal{J}}{\Gamma, y : B, x : A, \Delta \vdash \mathcal{J}} \\
 \text{(Contraction)} \frac{\Gamma, x : A, y : A, \Delta \vdash \mathcal{J}}{\Gamma, z : A, \Delta \vdash [z/x, z/y] \mathcal{J}}
 \end{array}$$

Data may be destroyed (by measurement).

Quantum data may *not* be duplicated (*No-cloning theorem*).

Computation Rules (based on [2])

We have the obvious β - and η -rules for $+$ and \otimes , plus:

$$\begin{aligned}
 &(\text{id}+) \text{inlr } (M, N) = M \\
 &(! + \text{id}) \text{inlr } (M, N) = N \\
 &\text{inlr } ((\text{id}+)M, (! + \text{id})M) = M \\
 &(\text{measure } \phi_1 \rightarrow M_1 \mid \dots \mid \phi_n \rightarrow M_n \mid 0 \rightarrow M_{n+1}) \\
 &\quad = \text{measure } \phi_1 \rightarrow M_1 \mid \dots \mid \phi_n \rightarrow M_n \\
 &(\text{measure } 1 \rightarrow M) = M \\
 &(\text{measure } \phi \otimes \psi \rightarrow M \mid \dots) = (\text{measure } \phi \rightarrow M \mid \psi \rightarrow M \mid \dots) \\
 &\quad \text{where } \phi \text{ and } \psi \text{ are scalars} \\
 &\text{let } x \otimes y = D(U, V)(M \otimes N) \text{ in } \begin{cases} \langle 0 | x \otimes \phi[y] \\ \langle 0 | XM \Leftrightarrow \langle 1 | M \end{cases} \\
 &\quad \Leftrightarrow \langle 0 | M \otimes \phi[UN] \\
 &(\text{let } x \otimes y = D(U, V)(M \otimes N) \text{ in } y) = \begin{pmatrix} \text{measure } \langle 0 | M \rightarrow UN \\ \langle 1 | M \rightarrow VN \end{pmatrix} \\
 &\quad \langle 0 | 0 \rangle \Leftrightarrow 1 \\
 &D(U, V)(|0\rangle \otimes N) = |0\rangle \otimes UN \\
 &\text{swap } (M \otimes N) = N \otimes M \\
 &IM = M \\
 &U(VM) = (UV)M \\
 &(U \otimes V)(M \otimes N) = UM \otimes VN
 \end{aligned}$$

where $D(U, V) = \begin{pmatrix} U & O \\ O & V \end{pmatrix}$.

Example — The Bell State

$$\text{CNOT} \stackrel{\text{def}}{=} D(I, X) \quad \text{Bell} \stackrel{\text{def}}{=} \text{CNOT}(H|0\rangle \otimes |0\rangle)$$

We can express the fact that the qubits in a Bell state are 7 entangled as follows:

$$\begin{aligned}
 &(\text{let } x \otimes y = \text{Bell} \text{ in } \langle 0 | x) \\
 &\Leftrightarrow (\text{let } x \otimes y = \text{CNOT}(H|0\rangle \otimes |0\rangle) \text{ in } \langle 0 | x \otimes \langle 0 | y) \\
 &\Leftrightarrow \langle 0 | H | 0 \rangle \otimes \langle 0 | 0 \rangle \\
 &\Leftrightarrow \langle 0 | H | 0 \rangle \otimes 1 \\
 &\Leftrightarrow \text{let } x \otimes y = \text{CNOT}(H|0\rangle \otimes |0\rangle) \text{ in } \langle 0 | x \otimes 1 \\
 &\Leftrightarrow \text{let } x \otimes y = \text{Bell} \text{ in } \langle 0 | x
 \end{aligned}$$

Bibliography

- Robin Adams. Qpel: Quantum programming and effect language. In *11th Workshop on Quantum Physics and Logic*, pages 133–153, 2014.
- Sam Staton. Algebraic effects, linearity, and quantum programming languages. In *POPL'15*, pages 395–406, 2015.