## RADBOUD UNIVERSITY NIJMEGEN

## EXAM Formal Languages, Grammars and Automata, NWI-MOL090

# June 25, 2013, 8.30 – 11.30, LIN 5

The maximum number of points per question is given in the margin. The last exercise (7) is "bonus" The grade is obtained by **dividing the number of points by** 10, rounding down to a 10 if needed.

1. Consider the following regular expressions

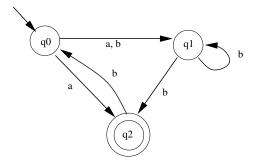
$$r_1 = a(ba)^*b$$
  $r_2 = (a(ba)^*b)^*$   $r_3 = (ab)^*ab$   $r_4 = (ab)^*$ 

- (5)
- (a) Give 2 regular expressions  $r_i$  and  $r_j$   $(i \neq j)$  such that  $L(r_i) = L(r_j)$ . Show that your answer is correct.
- (5) (b) Give 2 regular expressions  $r_i$  and  $r_j$  such that  $L(r_i) \neq L(r_j)$ . Show that your answer is correct.
  - 2. Consider the language

 $L := \{ w \in \{a, b\}^* \mid \#_a(w) \text{ is divisible by 3} \}$ 

NB "divisible by 3" is in Dutch: "deelbaar door 3"

- (6) (a) Give a regular expression e such that L(e) = L.
- (6)
- (b) Give a regular grammar G such that L(G) = L.
  - 3. Consider the following NFA (non-deterministic finite automaton) M



- (6) (a) Indicate for each of the following words whether they are accepted by M:
   *abba*, *ababa*, *abab*. Explain your answer.
- (7) (b) Give a regular expression e such that L(e) = L(M). Explain your answer.
- (7) (c) Construct a DFA M' that accepts the same language as M.

#### Continue on other side

4. Give a DFA (deterministic finite automaton) that accepts the language  $L_2$  where

 $L_2 = \{ w \in \{a, b\}^* \mid \#_a(w) \text{ is even and } w \text{ does not end with } aa \}.$ 

5. Consider the language over  $\{a, b, c\}$ 

(9)

 $L := \{ c \, w \, c \, v \, c \mid w, v \in \{a, b\}^* \text{ and } |w| = |v| \}$ 

So, w and v should be of equal length and not contain the symbol c.

- (8) (a) Give a context-free grammar G such that L = L(G).
- (8) (b) Give a pushdown automaton M such that L = L(M).
- (8) (c) Prove that L is not regular.
  - 6. Let M be the PDA with

$$Q = \{q_0, q_1\} \quad \delta(q_0, a, \lambda) = \{[q_0, A]\}$$
  

$$\Sigma = \{a, b, c\} \quad \delta(q_0, b, \lambda) = \{[q_0, B]\}$$
  

$$\Gamma = \{A, B\} \quad \delta(q_0, \lambda, A) = \{[q_1, \lambda]\}$$
  

$$F = \{q_1\} \quad \delta(q_1, c, A) = \{[q_1, \lambda]\}$$
  

$$\delta(q_1, \lambda, B) = \{[q_1, \lambda]\}$$

- (4) (a) Draw a state diagram for M.
- (6) (b) Check which of the following words is in L(M) and explain your answer: *bba*, *abba* and *baaacc*.
- (5) (c) Is  $L(b^*a) \subseteq L(M)$ ? Explain your answer.
- (5) (d) For which  $p, m, n \in \mathbb{N}$  do we have  $b^p a^m c^n \in L(M)$ ? Explain your answer.
- (5) (e) Describe L(M) using set-notation.
- (8) 7. [BONUS] Let  $\Sigma = \{a, b\}$ . Prove that, if L is regular over  $\Sigma$ , then Max(L) is also regular over  $\Sigma$ , where

$$Max(L) = \{ w \in L \mid \forall v \in \Sigma^* (w \, v \in L \to v = \lambda) \}$$

So, Max(L) consists of the words in L that cannot be 'extended'.

## END