

Huygens College Reflection

Assignment 7, Tuesday, Jan. 12, 2016

Remember the definitions of **true** := $\lambda xy.x(\equiv \mathbf{K})$ and **false** = $\lambda xy.y(=_{\beta} \mathbf{K}I)$, which are also written (on the slides of the course) as **T** := $\lambda xy.x$ and **F** := $\lambda xy.y$.

Exercise 1

(i) Show that there is a λ -term G that checks whether a term is an abstraction, that is:

$$\begin{aligned} G[x] &= \mathbf{F} \\ G[PQ] &= \mathbf{F} \\ G[\lambda x.P] &= \mathbf{T}. \end{aligned}$$

(See also the last exercise of the previous week.)

(ii) Show that there is a term R such that

$$\begin{aligned} R[M] &= \mathbf{T} \text{ if } M \text{ is a redex} \\ R[M] &= \mathbf{F} \text{ if } M \text{ is not a redex} \end{aligned}$$

(A redex is a term of the shape $(\lambda x.P)Q$.)

(iii) Show that there is a λ -term Norm that checks if a term is in *normal form*, so

$$\begin{aligned} \text{Norm}[M] &= \mathbf{T} \text{ if } M \text{ is in normal form} \\ \text{Norm}[M] &= \mathbf{F} \text{ if } M \text{ is not in normal form} \end{aligned}$$

(A term is *in normal form* if it contains no redex.)

(iv) You may need to define a conjunction $\&$ on Booleans first satisfying:

$$\begin{aligned} \& \mathbf{T} \mathbf{T} &= \mathbf{T} \\ \& \mathbf{T} \mathbf{F} &= \mathbf{F} \\ \& \mathbf{F} \mathbf{T} &= \mathbf{F} \\ \& \mathbf{F} \mathbf{F} &= \mathbf{F} \end{aligned}$$

Exercise 2

(i) Show that there is no term H' satisfying

$$\begin{aligned} H'[MN] &= \mathbf{T} \text{ if } MN \text{ has a normal form} \\ H'[MN] &= \mathbf{F} \text{ if } MN \text{ has no normal form} \end{aligned}$$

Hint: Show that, if H' exists, then we can also define the term B that solves the “blank tape” problem. (See course slides.)

(ii) Show that there is no term T satisfying

$$\begin{aligned} T \ulcorner M \urcorner &= \mathbf{T} \text{ if } M N \text{ has a normal form for all } N \\ T \ulcorner M \urcorner &= \mathbf{F} \text{ if } M N \text{ has no normal form for some } N \end{aligned}$$

Hint: Again, we can reduce B to T .

Exercise 3

(* Meaning: slightly more difficult) Show that there is a term C that contracts the *left-most redex* in a term M , so

$$C \ulcorner M \urcorner = \ulcorner N \urcorner$$

if N arises from M by contracting the left-most redex, if M contains a redex, and otherwise $N \equiv M$.