

Semantics and Domain theory

Exercises 7

1. Show that the untyped λ -term $\omega (= \lambda x.x x)$ is not typable in PCF. That is: show that there are no τ_1 and τ_2 such that $\vdash \mathbf{fn} x : \tau_1. x x : \tau_2$.
2. (a) Suppose that the term $\mathbf{mult} : \mathbf{nat} \rightarrow \mathbf{nat} \rightarrow \mathbf{nat}$ defines multiplication in PCF. Give a PCF term that defines the exponentiation function $\mathbf{exp} : \mathbf{nat} \rightarrow \mathbf{nat} \rightarrow \mathbf{nat}$. (So $\mathbf{exp} n m$ should denote n^m .)
 (b) Let $p : \mathbf{nat} \rightarrow \mathbf{nat}$. Define a term $N : \mathbf{nat}$ that denotes the smallest number n such that $p(n) = 0$ and $\forall i < n (p(i) > 0)$.
3. Show that, if $Q \Downarrow_{\tau} V$, then $(\mathbf{fn} x : \tau. \mathbf{fn} y : \tau. y) P Q \Downarrow V$. (NB. V denotes an arbitrary *value*.)
4. To prove that PCF evaluation is deterministic, we prove (in Proposition 5.4.1) that the following set is closed under the rules of Fig.3

$$\{(M, \tau, V) \mid M \Downarrow_{\tau} V \wedge \forall V' (M \Downarrow_{\tau} V' \Rightarrow V = V')\}$$

Show this for the cases of the rules $(\Downarrow_{\mathbf{if}1})$ and $(\Downarrow_{\mathbf{cbn}})$.

(An alternative way of looking at this is to prove the following:

$$M \Downarrow_{\tau} V \Rightarrow \forall V' (M \Downarrow_{\tau} V' \Rightarrow V = V')$$

by induction on the derivation of $M \Downarrow_{\tau} V$. Do only the cases when the last applied rule is $(\Downarrow_{\mathbf{if}1})$ or $(\Downarrow_{\mathbf{cbn}})$.)

5. Given the definition of plus (Exercise 5.6.3.)

$$\begin{aligned} \mathbf{plus} = & \mathbf{fix}(\mathbf{fn} p : \mathbf{nat} \rightarrow \mathbf{nat} \rightarrow \mathbf{nat}. \mathbf{fn} x : \mathbf{nat}. \mathbf{fn} y : \mathbf{nat}. \\ & \mathbf{if} \mathbf{zero}(y) \mathbf{then} x \mathbf{else} \mathbf{succ}(p x \mathbf{pred}(y))) \end{aligned}$$

Prove (by induction) that

$$\forall m, n (\mathbf{plus} \mathbf{succ}^m(0) \mathbf{succ}^n(0) \Downarrow_{\mathbf{nat}} \mathbf{succ}^{m+n}(0))$$

NB. First identify the proper statement that you need and that you can prove relatively easily by induction.