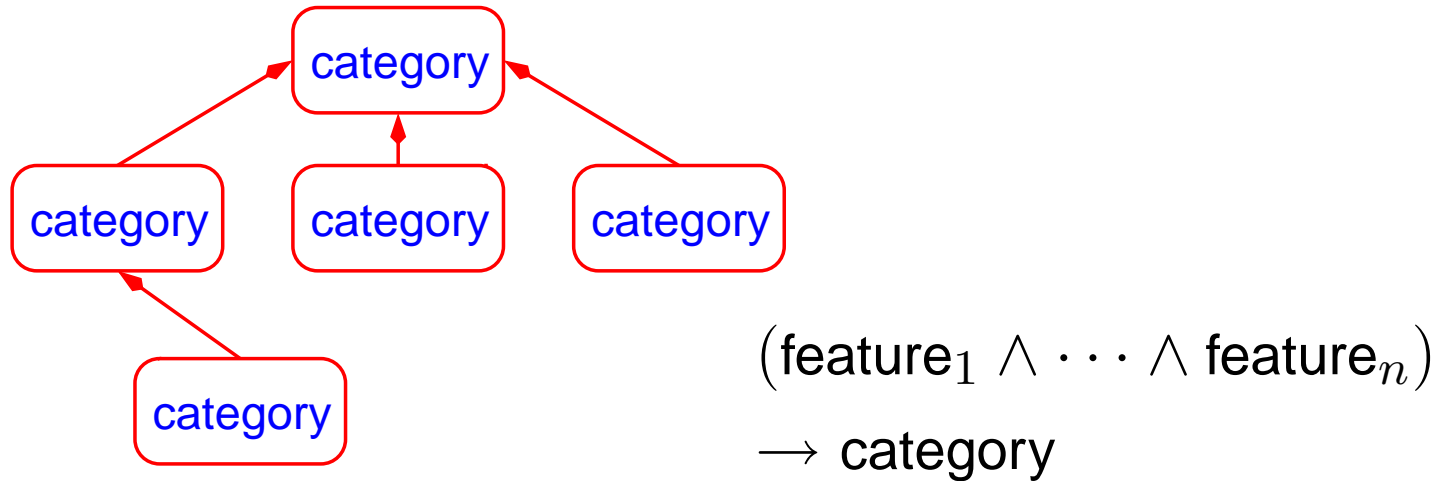


Model-based Reasoning

- 'Traditional' knowledge systems:
 - **Rule based** (heuristic rules)
 - if conditions then actions/conclusions fi
 - ⋮
 - Reasoning:
 - forward chaining: reasoning from facts to conclusions
 - backward chaining: reasoning from goals to facts
 - Recent: **business rules**
- Model-based systems: reasoning with understandable **model**, i.e., they have **intuitive** semantics

Heuristic rules and their disadvantages

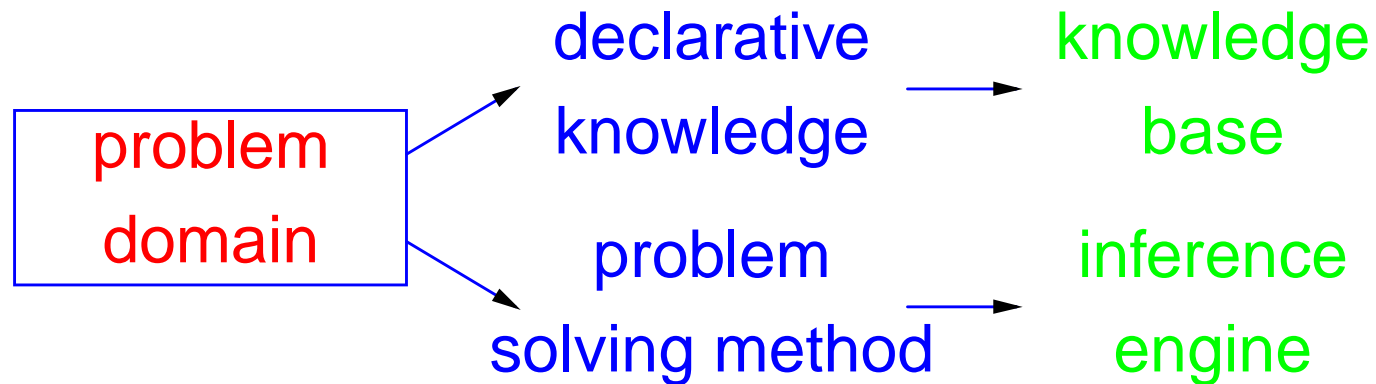
Problem solving based on heuristic rules:



Disadvantages:

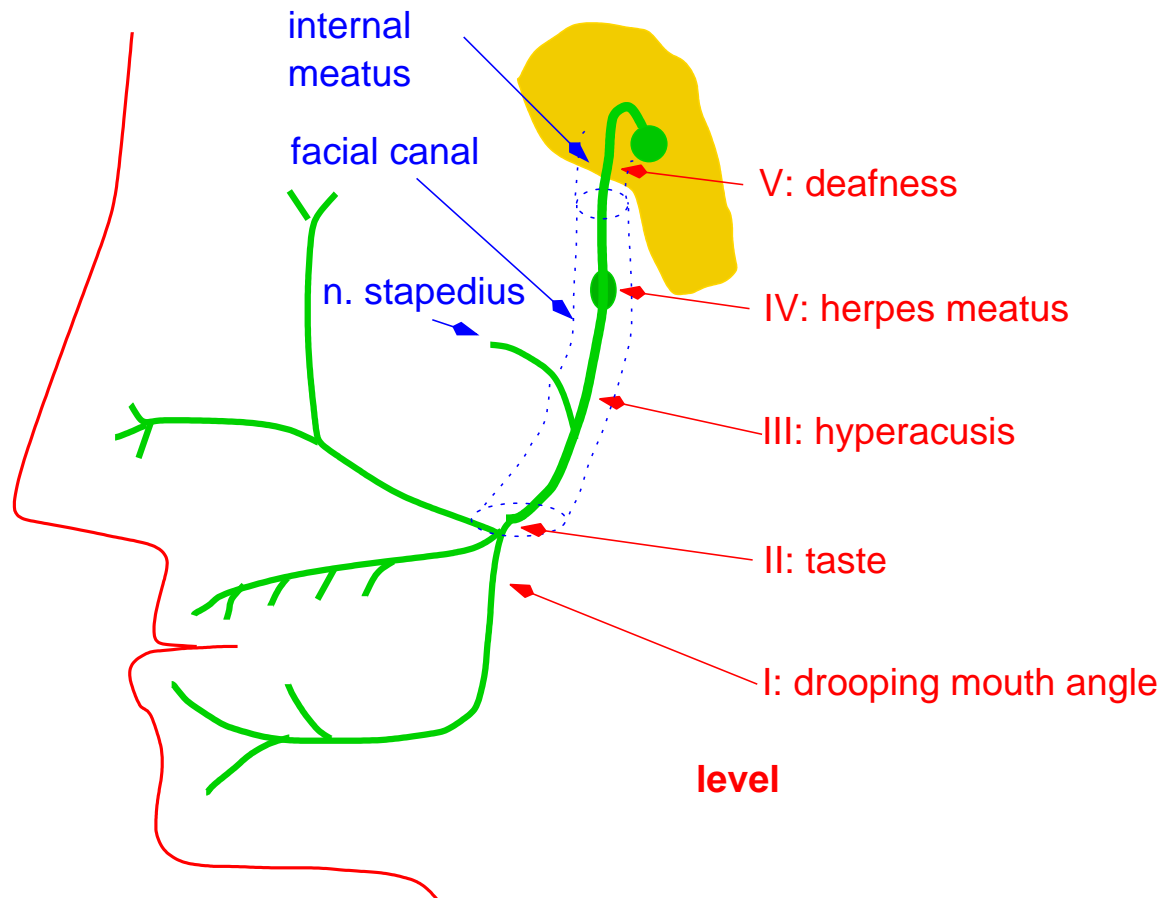
- no use of knowledge about structure and workings
- knowledge maintenance and updating is hard

Model use



- **Models:**
 - usually designed for handling **multiple** problems ⇒ **reuse**
 - capture instantaneous behaviour, temporal behaviour, structure
- **Methods:** diagnosis, decision making, prediction, planning

Medical diagnosis of facial palsy



Diagnosis:

$$\text{Symptoms-level}_{i+1} = \text{Symptoms-level}_i \cup \text{New-symptoms}$$

Drilling Automation for Mars

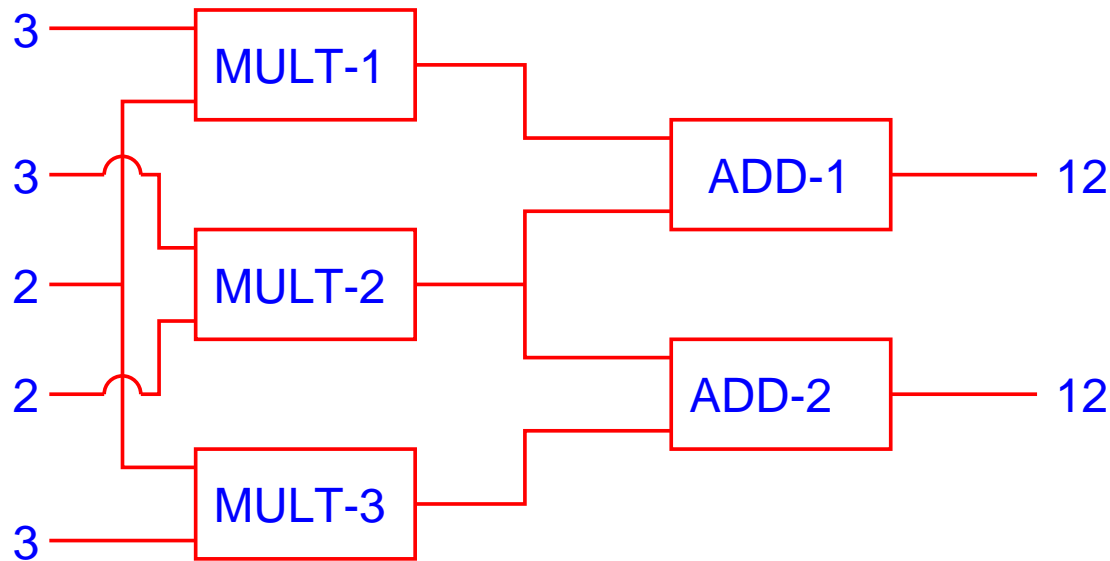


Autonomous drill
with sensors



Fault diagnosis

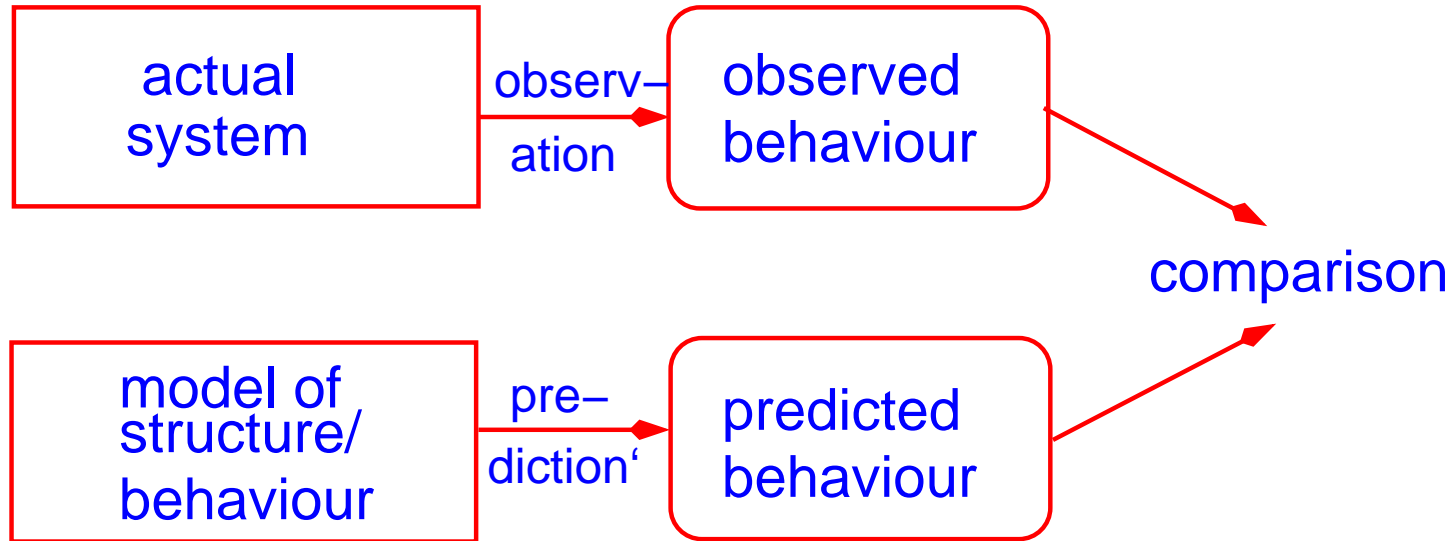
Structure and behaviour



Multiplier-adder

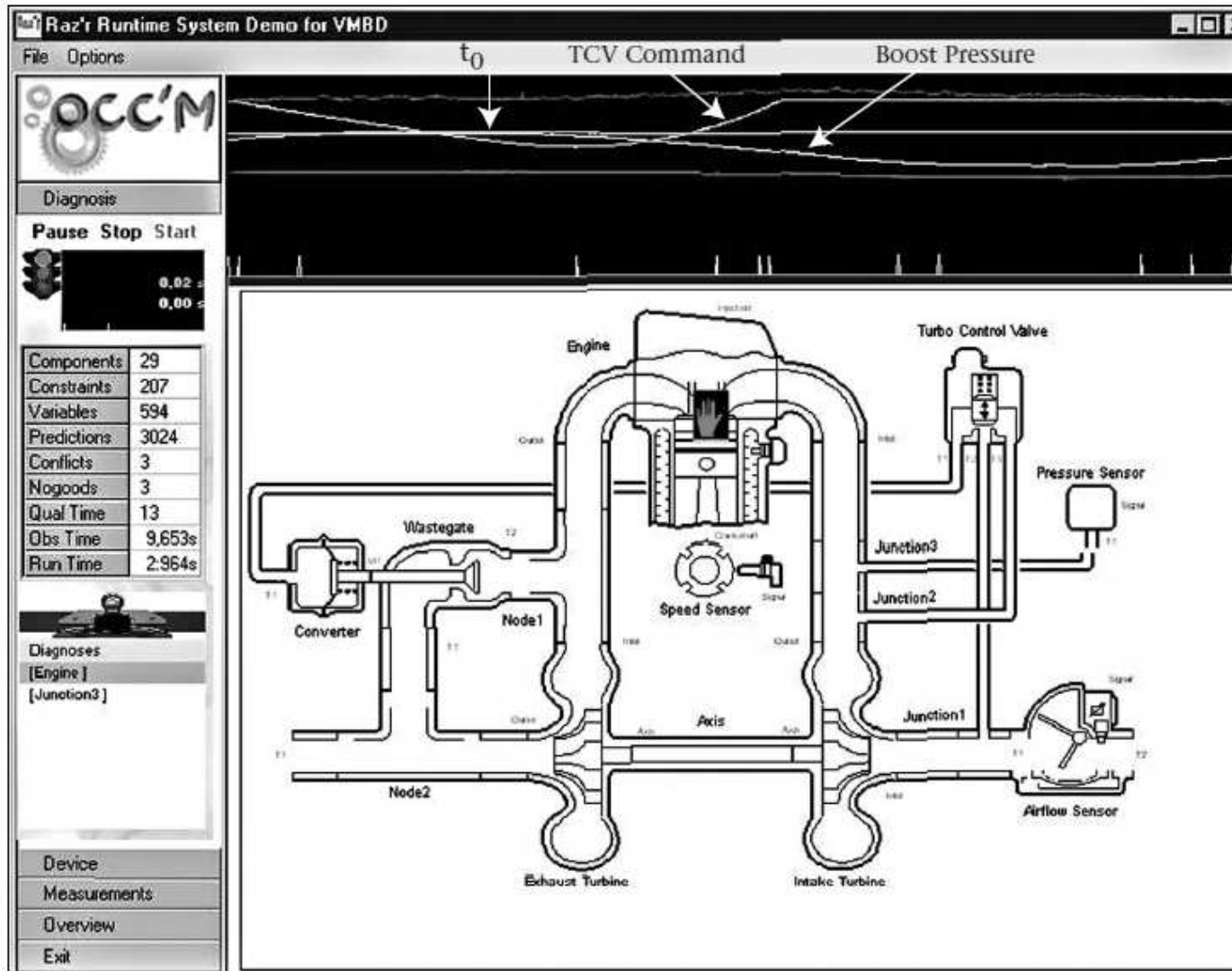
- **Structure:**
 - components: MULT-1, MULT-2, ...
 - wiring
- **Behaviour:**
 - behaviour of individual components
 - combined behaviour

Method: model-based diagnosis

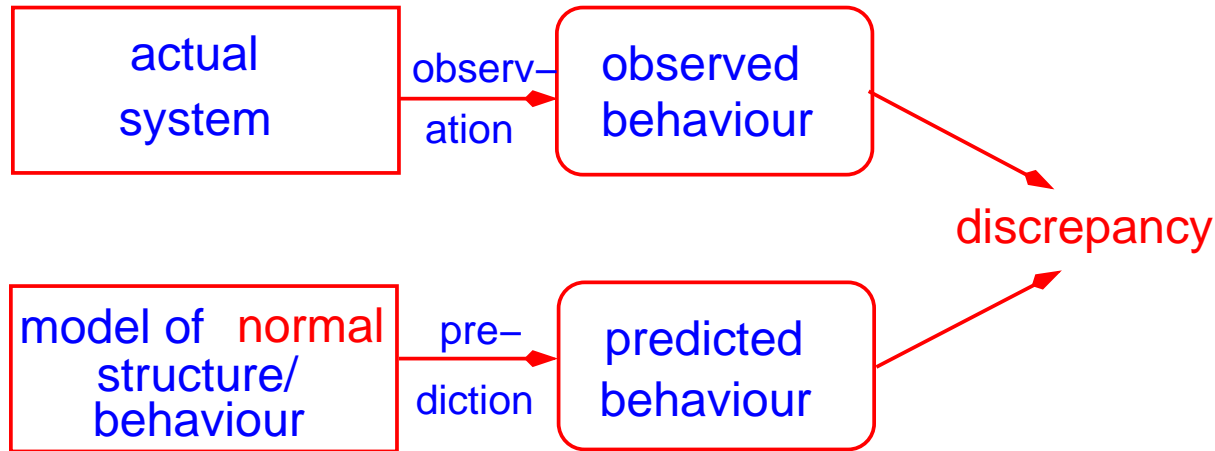


- Model: representation of **normal** or **abnormal** behaviour and, possibly, internal **structure**
- Formalisation:
 - *consistency-based diagnosis*, and
 - *abductive diagnosis*

OCC'M



Consistency-based diagnosis

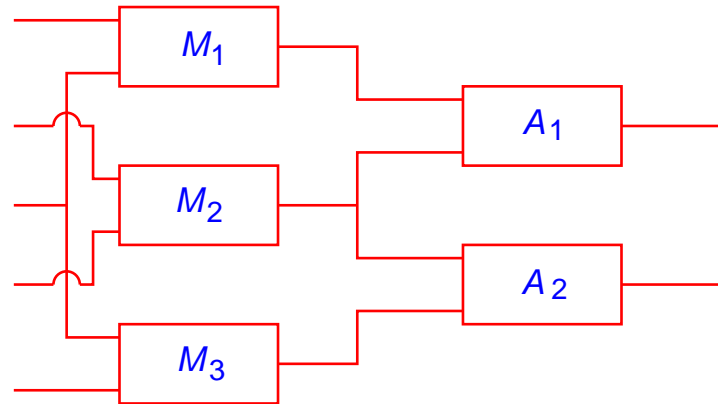


Difference between predicted behaviour and observed behaviour \Rightarrow **defect!**

Originators:

- R. Reiter, “A Theory of diagnosis from first principles”, *Artificial Intelligence*, vol. 32, 57–95, 1987.
- J. de Kleer, A.K. Macworth, and R. Reiter, “Characterising diagnoses and systems”, *Artificial Intelligence*, vol. 52, 197–222, 1992.

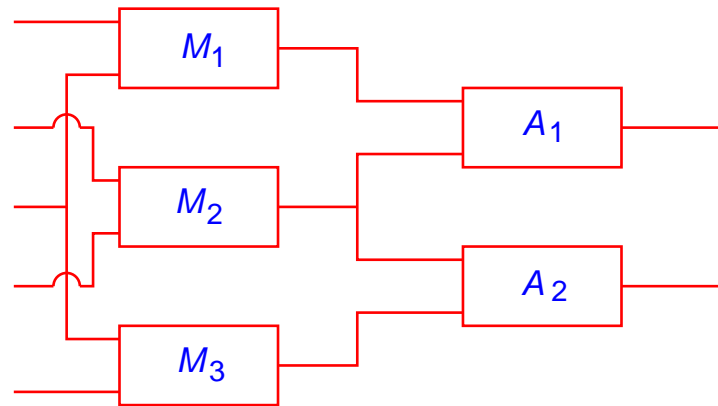
Normal behaviour



SYStem specification $SYS = (SD, COMPS)$:

- Components that **may** be **defective** (faulty):
 $COMPS = \{M_1, M_2, M_3, A_1, A_2\}$
- SD (**System Description**):
 - **generic** description of component behaviour (what the component does)
 - declaration of **components**: $MUL(M_1)$, $ADD(A_1)$
 - **connection** between components

Normal behaviour: formal



SYStem specification $\text{SYS} = (\text{SD}, \text{COMPS})$:

- **SD (System Description):**

$$\forall x (\text{MUL}(x) \rightarrow \text{in}_1(x) \times \text{in}_2(x) = \text{out}(x))$$

$$\forall x (\text{ADD}(x) \rightarrow \text{in}_1(x) + \text{in}_2(x) = \text{out}(x))$$

$$\text{MUL}(M_1), \text{MUL}(M_2), \text{MUL}(M_3), \text{ADD}(A_1), \text{ADD}(A_2)$$

$$\text{in}_1(A_1) = \text{out}(M_1), \text{in}_2(A_1) = \text{out}(M_2)$$

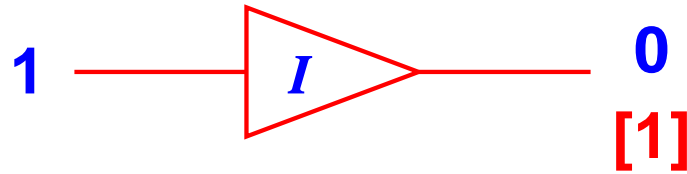
$$\text{in}_1(A_2) = \text{out}(M_2), \text{in}_2(A_2) = \text{out}(M_3)$$

- **COMPS** = $\{M_1, M_2, M_3, A_1, A_2\}$

Ab predicate

- $Ab(c)$: component c is *abnormal*
- $\neg Ab(c)$: component c is not abnormal, i.e. *normal*

Example (Inverter I):



$SD = \{\forall x((INV(x) \wedge \neg Ab(x)) \rightarrow \neg(\text{out}(x) = \text{in}(x))), INV(I)\}$

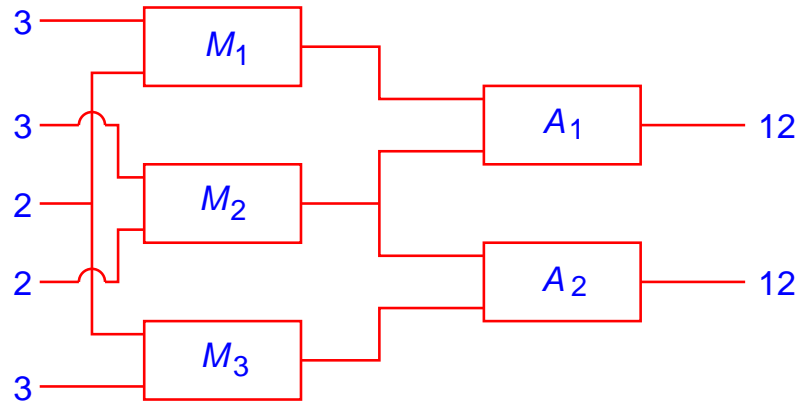
- **Input:** $\text{in}(I) = 1$; **observed output:** $\text{out}(I) = 1$

$SD \cup \{\text{in}(I) = 1, \text{out}(I) = 1\} \cup \{\neg Ab(I)\} \models \perp$

$SD \cup \{\text{in}(I) = 1, \text{out}(I) = 1\} \cup \{Ab(I)\} \not\models \perp$

(assumption that I is (ab)normal is (in)consistent)

Normal behaviour formal



SYStem specification $\text{SYS} = (\text{SD}, \text{COMPS})$:

- **SD (System Description):**

$$\forall x((\text{MUL}(x) \wedge \neg \text{Ab}(x)) \rightarrow \text{in}_1(x) \times \text{in}_2(x) = \text{out}(x))$$

$$\forall x((\text{ADD}(x) \wedge \neg \text{Ab}(x)) \rightarrow \text{in}_1(x) + \text{in}_2(x) = \text{out}(x))$$

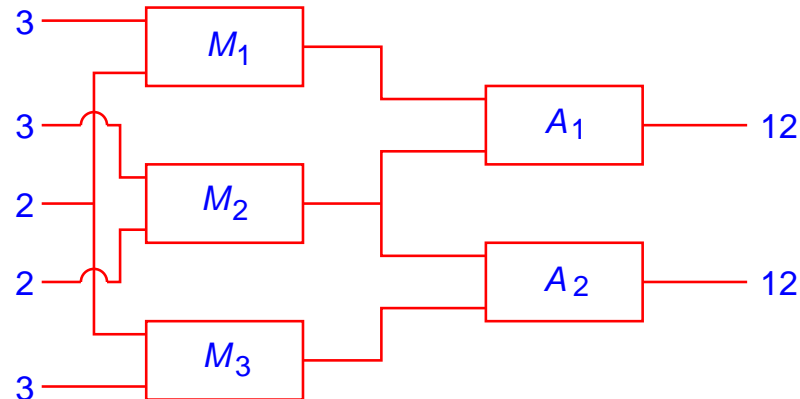
$$\text{MUL}(M_1), \text{MUL}(M_2), \text{MUL}(M_3), \text{ADD}(A_1), \text{ADD}(A_2)$$

$$\text{in}_1(A_1) = \text{out}(M_1), \text{in}_2(A_1) = \text{out}(M_2)$$

$$\text{in}_1(A_2) = \text{out}(M_2), \text{in}_2(A_2) = \text{out}(M_3)$$

- **COMPS** = $\{M_1, M_2, M_3, A_1, A_2\}$

Prediction of normal behaviour



- System specification $SYS = (SD, COMPS)$

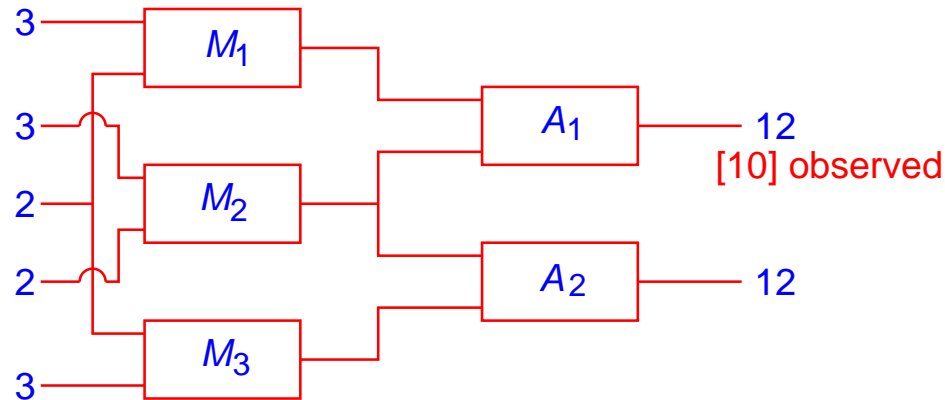
- A **prediction**:

$SD \cup \{\neg Ab(c) \mid c \in COMPS \text{ is not defective}\} \cup \text{Inputs} \models$
Behaviour (Inputs stands for all observations)

- Example:

$SD \cup \{\neg Ab(M_1), \neg Ab(M_2), \neg Ab(A_1)\} \cup \{\text{in}_1(M_1) = 3,$
 $\text{in}_2(M_1) = 2, \text{in}_1(M_2) = 3, \text{in}_2(M_2) = 2\} \models \text{out}(A_1) = 12$

There is a fault!



Let $(\text{out}(A_1) = 10) \in \text{Inputs}$, then:

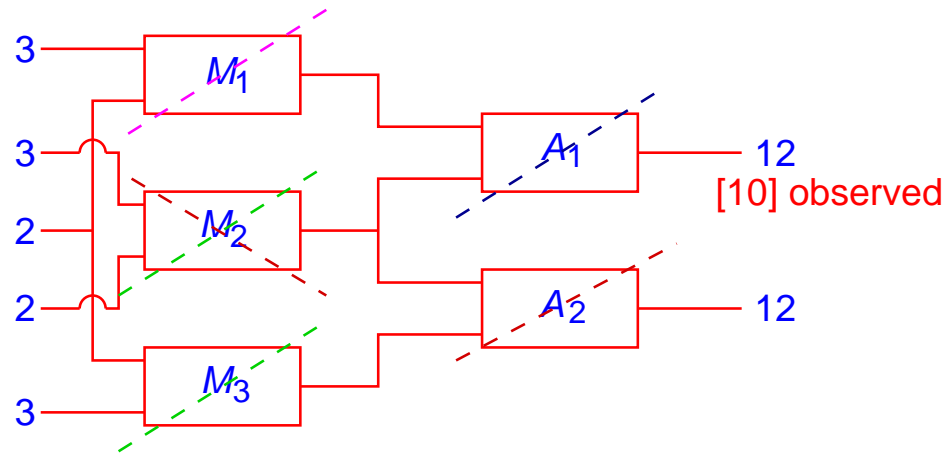
$$\text{SD} \cup \{\neg \text{Ab}(c) \mid c \in \text{COMPS}\} \cup \text{Inputs} \models \perp$$

because:

- $\text{SD} \cup \{\neg \text{Ab}(c) \mid c \in \text{COMPS}\} \cup \text{Inputs} \models \text{out}(A_1) = 12$, and
- $\text{SD} \cup \{\neg \text{Ab}(c) \mid c \in \text{COMPS}\} \cup \text{Inputs} \models \text{out}(A_1) = 10$

\Rightarrow **faulty component**

Which components are faulty?



Possible **diagnoses** (faulty components) D :

- $D = \{A_1\}, \{M_1\}, \{M_2, M_3\}, \{A_2, M_2\}$, because
$$\text{SD} \cup \{\neg \text{Ab}(c) \mid c \in \text{COMPS} - D\} \\ \cup \{\text{Ab}(c) \mid c \in D\} \cup \text{Inputs} \neq \perp$$
- D must be the *smallest* set because, $D = \text{COMPS}$ would also be a diagnosis otherwise

\Rightarrow **multiple diagnoses**

Diagnostic problem

- System specification $SYS = (SD, COMPS)$
- Diagnostic problem $DP = (SYS, OBS)$, with OBS a set of observations
- A diagnosis D : smallest (subset minimal) set of components, such that

$$SD \cup OBS \cup \{Ab(c) \mid c \in D\} \cup \{\neg Ab(c) \mid c \in COMPS - D\}$$

is consistent

Example:

$$OBS = \{in_1(M_1) = 3, in_2(M_1) = 2, in_1(M_2) = 3, in_2(M_2) = 2, \\ in_1(M_3) = 2, in_2(M_3) = 3, out(A_1) = 10, out(A_2) = 12\}$$

Algorithms

- **Enumerate** all diagnoses: #P complete (NP hard for enumeration):

$$Ab(A_1) \wedge \neg Ab(A_2) \wedge \neg Ab(M_1) \wedge \neg Ab(M_2) \wedge \neg Ab(M_3)$$

$$Ab(A_1) \wedge Ab(A_2) \wedge \neg Ab(M_1) \wedge \neg Ab(M_2) \wedge \neg Ab(M_3)$$

$$Ab(A_1) \wedge Ab(A_2) \wedge Ab(M_1) \wedge \neg Ab(M_2) \wedge \neg Ab(M_3)$$

⋮

- **Heuristic methods:**
 - hitting set algorithm (Reiter)
 - assumption-based truth maintenance system (ATMS, De Kleer)
- **Restrictions:** for example, only maximally 2 defects, then complexity upperbound

Basic problem: which idea should underly such algorithms?

Conflict set

Let $CS \subseteq COMPS$ be a set of components, then **CS** is called a **conflict set** iff

$$SD \cup OBS \cup \{\neg Ab(c) \mid c \in \mathbf{CS}\}$$

is inconsistent ($SD \cup OBS \cup \{\neg Ab(c) \mid c \in CS\} \models \perp$)

Proposition: For each $D \subseteq COMPS$ that is a **diagnosis** and each conflict set CS it holds that: $D \cap CS \neq \emptyset$

Proof: $SD \cup OBS \cup \{\neg Ab(c) \mid c \in COMPS - D\} \not\models \perp$, with D subset minimal $\Rightarrow COMPS - D$ is subset maximal, hence

$$SD \cup OBS \cup \underbrace{\{\neg Ab(c) \mid c \in COMPS - D\} \cup \{\neg Ab(c')\}}_{CS} \models \perp$$

for c' in CS

Basic ideas: hitting sets

- Determine conflict sets of the diagnostic problem DP
- Each conflict set CS has at least one element in common with a diagnosis D :

$$D = \{c_1, c_2, \dots, c_m\}$$

/ \

$$CS_1 = \{\dots, c_1, \dots\} \quad CS_m = \{\dots, c_m, \dots\}$$

- Compute so-called **hitting sets**:
 - Let F be a set of sets, and
 - $H \subseteq \bigcup_{S \in F} S$,
 - H is a **hitting set** if for all $S \in F : H \cap S \neq \emptyset$
- Example: $F = \{\{1, 2\}, \{3, 4\}\}$, then $H = \{2, 4\}$ is a (non-unique) hitting set ($H = \{1, 3\}$ is also a hitting set)

Diagnosis as hitting set

Theorem: D is a diagnosis for diagnostic problem $DP = (SYS, OBS)$ iff D is a minimal hitting set for all conflict sets of DP

Proof (sketch):

- Prop. page 19: for all conflict sets CS : $CS \cap D \neq \emptyset$.
Thus, D is a hitting set
- D is also a *minimal* hitting set, as $COMPS - D$ is no conflict set, whereas $\{c\} \cup (COMPS - D)$ is a conflict set for any $c \in D$

Hitting-set tree

We construct a **tree structure** for computing diagnoses

Some definitions:

- Let F be a set of sets
- Let $T = (V, E, l_V, l_E)$ be a **labelled tree**, with V a set of **nodes** and $E \subseteq V \times V$ a set of **edges**, and
 - l_V a **node label function**:

$$l_V : V \rightarrow F \cup \{\checkmark\}$$

- l_E an **edge label function**:

$$l_E : E \rightarrow \bigcup_{S \in F} S$$

Node label function

Let F be a set of sets:

- l_V a node label function:

$$l_V : V \rightarrow F \cup \{\checkmark\}$$

with

$$l_V(v) = \begin{cases} S & \text{if } S \in F, S \neq \emptyset \\ \checkmark & \text{otherwise} \end{cases}$$

- Example: $F = \{\{1, 2\}, \{4, 5\}\}$
 - Nodes $V = \{u, v, w\}$
 - $l_V(u) = \{1, 2\}$, $l_V(v) = \{4, 5\}$, and $l_V(w) = \checkmark$

Edge label function

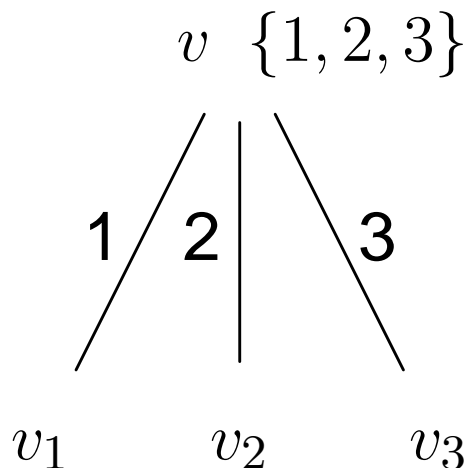
Let F be a set of sets:

- l_E an edge label function:

$$l_E : E \rightarrow \bigcup_{S \in F} S$$

with if $l_V(v) = S$ and $\forall s \in S: l_E(v, v_s) = s$

Example:



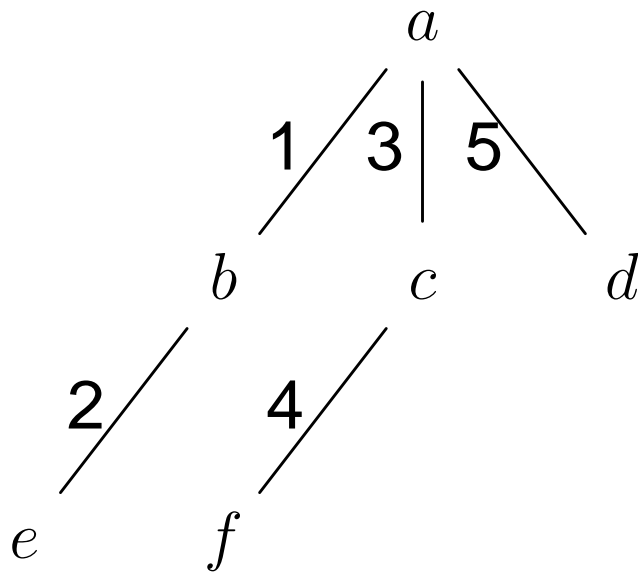
- $F = \{\{4, 5\}, \{1, 2, 3\}\}$, and $S = \{1, 2, 3\}$
- $l_V(v) = \{1, 2, 3\}$
- $l_E(v, v_1) = 1, l_E(v, v_2) = 2, l_E(v, v_3) = 3$

Construction of hitting sets

Let $T = (V, E, l_V, l_E)$ be a labelled tree, then the **hitting set** $H(v)$ for node v is defined as:

$$H(v) = \{l_E(u, w) \mid (u, w) \text{ is on the path from the root to } v\}$$

Example:



- $H(a) = \emptyset$
- $H(b) = \{1\}$
- $H(e) = \{1, 2\}$
- $H(f) = \{3, 4\}$

Hitting-set algorithm

- F is the set of conflict sets, which is **initially empty**
- Let node v_s be a child of node v , then $l_V(v_s) = \text{CS}$ if there exists a $\text{CS} \in F$ with

$$\text{CS} \cap H(v_s) = \emptyset$$

(CS is not yet covered by $H(v_s)$ and we have to extend the path)

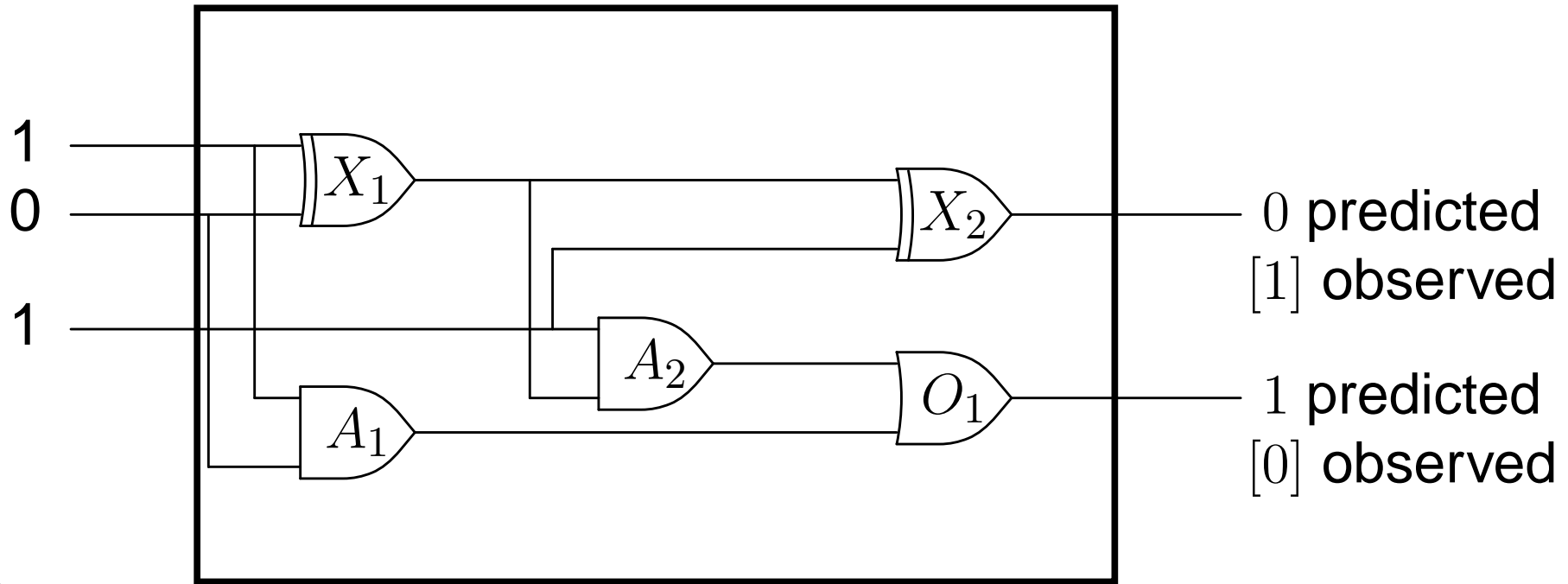
- If no suitable $\text{CS} \in F$, call **logical reasoning program TP**:
 - Call: $\text{TP}(\text{SD}, \text{COMPS} - H(v_s), \text{OBS})$
 - Returns: conflict set CS if $\text{SD} \cup \text{OBS} \cup \{\neg \text{Ab}(c) \mid c \in \text{COMPS} - H(v_s)\} \models \perp$, with $\text{CS} \cap H(v_s) = \emptyset$ and $\text{CS} \subseteq \text{COMPS} - H(v_s)$ otherwise, \checkmark (consistent)

Hitting-set algorithm

```
Diagnose(SD, COMPS, OBS)
{
  generate HS tree by calling
    CS <- TP(SD, COMPS - H(v), OBS);
  (F is build op from these CS's)
  leaves v with ✓
    determine diagnosis H(v)
  determine subset-minimal H(v)
    in the HS tree
}
```

```
TP(A, B, C)
{
  use resolution on A U B U C
}
```

Example: full-adder



SD:

$$\forall x((\text{ANDG}(x) \wedge \neg \text{Ab}(x)) \rightarrow (\text{out}(x) = \text{in}_1(x) \wedge \text{in}_2(x)))$$

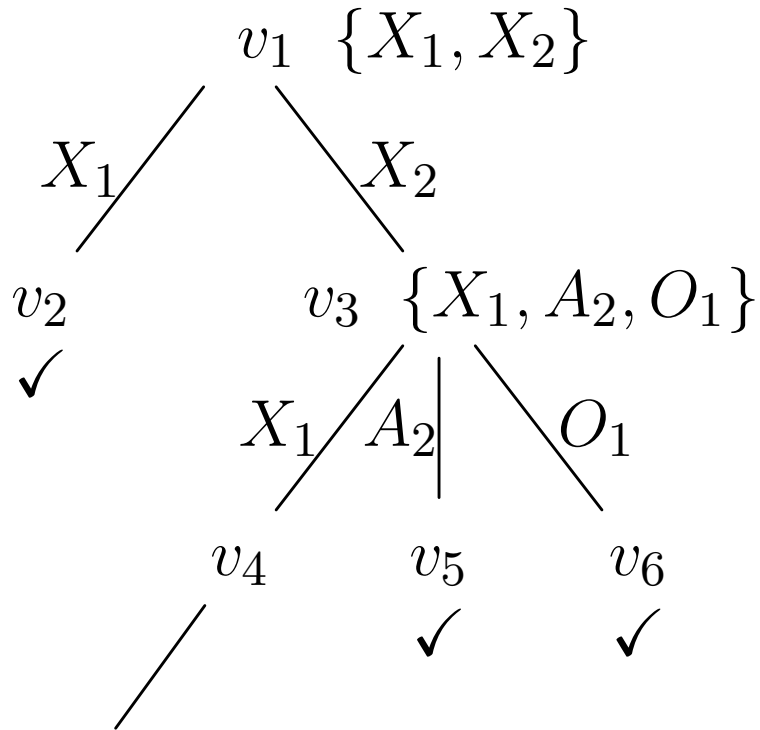
$$\forall x((\text{ORG}(x) \wedge \neg \text{Ab}(x)) \rightarrow (\text{out}(x) = \text{in}_1(x) \vee \text{in}_2(x)))$$

⋮

$$\text{ORG}(O_1), \text{ANDG}(A_1), \text{XORG}(X_1), \dots$$

$$\text{COMPS} = \{A_1, A_2, X_1, X_2, O_1\}$$

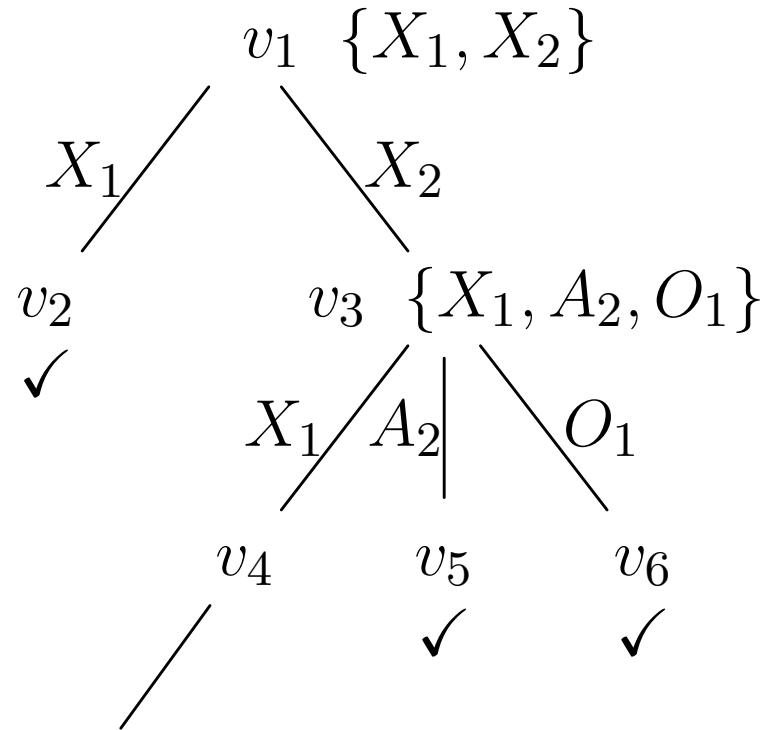
Example HS tree



1. $CS_1 \leftarrow TP(SD, COMPS, OBS);$
 $CS_1 \leftarrow \{X_1, X_2\}$
2. $CS_2 \leftarrow TP(SD, COMPS - \{X_1\}, OBS);$
 $CS_2 \leftarrow \checkmark$ (diagnosis found)
3. $CS_3 \leftarrow TP(SD, COMPS - \{X_2\}, OBS);$
 $CS_3 \leftarrow \{X_1, A_2, O_1\}$
4. \vdots

Diagnoses D : $\{X_1\}$, $\{X_2, A_2\}$, $\{X_2, O_1\}$
 (note that $\{X_2, X_1\}$ not subset minimal)

Pruning of the HS tree



Note that there is no need to extend the hitting set $H(v_4)$ (as $\{X_2, X_1\}$ is not subset minimal)

\Rightarrow pruning of the hitting-set tree