## Is $\chi_{n}$ a power function?

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13 July 2023

## Introduction

## Introduction to $\chi_{n}$

Definition $1\left(\chi_{n}\right)$
The map $\chi_{n}: \mathbb{F}_{2}^{n} \rightarrow \mathbb{F}_{2}^{n}, x \mapsto y$ is given by:

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## Univariate polynomials

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## Definition 2 (Normal basis)

Consider $\mathbb{F}_{2} \subset \mathbb{F}_{2^{n}}$. Then $\beta \in \mathbb{F}_{2^{n}}$ is called a normal element of $\mathbb{F}_{2^{n}}$ if the set $\left\{\beta, \beta^{2}, \beta^{2^{2}}, \ldots, \beta^{2^{n-1}}\right\}$ is a linear independent set. This set is then called a normal basis of $\mathbb{F}_{2^{n}}$.

## Univariate polynomials

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\begin{aligned}
& \mathbb{F}_{2}^{n} \xrightarrow{\chi_{n}} \mathbb{F}_{2}^{n} \quad \phi(\vec{x})=\sum_{i=0}^{n-1} x_{i} \beta^{q^{i}} \\
& \phi \mid \\
& \mathbb{F}_{2^{n}} \xrightarrow[\widehat{\widehat{\chi_{n}}}]{ } \mathbb{F}_{2^{n}}
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## Theorem 3

If $F: \mathbb{F}_{2}^{n} \rightarrow \mathbb{F}_{2}^{n}$ is shift invariant and the isomorphism $\phi$ is induced by a normal element, then $\widehat{F}$ has coefficients in $\mathbb{F}_{2}$.

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\begin{aligned}
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& \left.\phi\right|_{\downarrow} \\
& \mathbb{H}_{2^{n}} \xrightarrow[\widehat{\widehat{\chi n}}]{ } \mathbb{F}_{2^{n}}
\end{aligned}
$$

## Example 2

Consider the map $\chi_{3}$. Let $\mathbb{F}_{2^{3}}:=\mathbb{F}_{2}(\alpha)=\mathbb{F}_{2}[X] /\left(X^{3}+X+1\right)$. Then $\alpha^{3}$ is a normal element. We define $\widehat{\chi_{3}}:=\phi \circ \chi_{3} \circ \phi^{-1}$. By using Lagrange interpolation we find that $\widehat{\chi_{3}}(t)=t^{6}$.

## Power functions

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A power function is a polynomial function that can be represented by a single monomial in $\mathbb{F}_{2^{n}}[X]$. We write $*^{e}: \mathbb{F}_{2^{n}} \rightarrow \mathbb{F}_{2^{n}}$ for a power function, here $e \geq 0$.

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Why power functions?

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## Proposition 1

For any even $n$, there is no (normal) basis representation such that $\widehat{\chi_{n}}$ is a power function.

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## Proposition 1

For any even $n$, there is no (normal) basis representation such that $\widehat{\chi_{n}}$ is a power function.

## Proof.

Suppose that it does exist. Since $\chi_{n}\left((01)^{n / 2}\right)=0^{n}$, there needs to exist some nonzero $\alpha \in \mathbb{F}_{2^{n}}$ with $\alpha^{s}=0$ for some integer $s$.

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No! (For $n \neq 1,3$.)

## Proposition 1 (Excluding Mersenne-exponents)

If $n>3$ is such that $2^{n}-1$ is a prime number, then there exists no (normal) basis representation of $\chi_{n}$ such that $\widehat{\chi_{n}}$ is a power function.

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## Proof.

Since $2^{n}-1$ is a prime number, $\varphi\left(2^{n}-1\right)=2^{n}-2$. The order of $\chi_{n}$ is divisible by 4 for all $n>3$. The expression $2^{n}-2$ has only one factor 2 .

## Historical attempts

## State diagrams

## Definition 4 (State diagram)

Let $S$ be a set. The state diagram for a map $F: S \rightarrow S$ is a directed graph $(V, A)$, where $V=S$ and $A=\{(a, F(a)) \mid a \in S\}$.

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## Theorem 5 (Ahmad's Theorem)

Let $m, q$ be positive integers with $q=p^{n}$ for some prime number $p$ and $n \geq 1$. Let $*^{e}: \mathbb{F}_{q}^{*} \rightarrow \mathbb{F}_{q}^{*}, x \mapsto x^{e}$ be a power function. Then $*^{e}$ has a cycle of length precisely $m$ if and only if there exists some $t \mid q-1$ such that the order of $e$ modulo $t$ is equal to $m$.

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Not every length necessarily occurs!

## Corollary

Theorem 6 (Necessary conditions for $\chi_{n}$ to be a power function)
Let $n>3$ be an odd integer. Write $o:=\operatorname{ord}\left(\chi_{n}\right)=2^{\lfloor\lg (n)\rfloor}$. Then $\chi_{n}$ can only be a power function if $2^{n}-1$ factors as

$$
2^{n}-1=p_{1}^{e_{1}} \cdots p_{r}^{e_{r}},
$$

such that there exists some permutation $\sigma \in S_{r}$ with

$$
\begin{gathered}
\varphi\left(p_{\sigma(1)}^{e_{\sigma(1)}}\right) \text { is a multiple of } o \\
\varphi\left(p_{\sigma(2)}^{e_{\sigma(2)}}\right) \text { is a multiple of } \frac{o}{2} \\
\vdots \\
\varphi\left(p_{\sigma(t)}^{e_{\sigma(t)}}\right) \text { is a multiple of } 2
\end{gathered}
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for some $t<r$.

## Results

Using these conditions, we can verify ${ }^{1}$ that $\chi_{n}$ is not a power function for any $n \leq 1115$, except for $n=63$ and $n=441$.

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Remaining cases:

- $n=63$ :
- $\approx 2^{12.59}$ out of $\approx 2^{62.742}$ possible $e$;

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- $2^{35.322}$ out of $\approx 2^{440.742}$ possible $e$;

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## Proof technique

## Differential distributions

## Definition 7 (Differential probability (Biham, Shamir))

Let $f: G \rightarrow H$ be a map between finite groups $G$ and $H$. Let $g \in G$ and $h \in H$ be arbitrary. Then we define the differential probability of $f$ at $(g, h)$ as

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## Differential distribution for $\chi$

## Proposition 2 (Differential probabilities for $\chi$ (Daemen))

Let $n>1$ be an arbitrary odd integer. Let $a \in \mathbb{F}_{2}^{n}$ be arbitrary. Then for any compatible $b \in \mathbb{F}_{2}^{n}$ we have $\mathrm{DP}_{\chi_{n}}(a, b)=2^{-w(a)}$, where

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w(a)= \begin{cases}n-1 & \text { if } a=1^{n} ; \\ \operatorname{wt}(a)+r & \text { else, }\end{cases}
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where $r$ is the number of (cyclic) 001-substrings in $a$.

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- $a^{\prime}=10^{n-1} \Longrightarrow \mathrm{DP}_{\chi_{n}}\left(a^{\prime}, b\right)=\frac{1}{4}$.


## Invariant

Proposition 3 (Differential probabilities under linear isomorphisms)
Let $G \xlongequal{\varrho} H$ be isomorphic groups. Let $f: G \rightarrow G$ be a map and let $\hat{f}: H \rightarrow H$ be the map induced through the isomorphism. Then $\mathrm{DP}_{\hat{f}}(g, h)=\mathrm{DP}_{f}\left(\varphi^{-1}(g), \varphi^{-1}(h)\right)$ for all $g, h \in H$.

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Let $0 \leq e \leq 2^{n}-1$ and let $f=*^{e}: \mathbb{F}_{2^{n}} \rightarrow \mathbb{F}_{2^{n}}$ be a power function. Then $\mathrm{DP}_{f}(a, b)=\mathrm{DP}_{f}\left(y a, y^{e} b\right)$ for all $y \in \mathbb{F}_{2^{n}}^{*}$.

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Thus, we have that the rows of the DDT all have the same number of occurrences of $0,2,4, \ldots$.

## Conclusion and corollary

## Theorem 8

Let $n \neq 1,3$ be a positive integer. Then $\widehat{\chi_{n}}$ is not a power function.

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There is no function $F_{n}$ that is extended affine equivalent to $\chi_{n}$ $\left(A F_{n} B+C=\chi_{n}\right)$, such that $\widehat{F_{n}}$ is a power function.

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Thank you for your attention!


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